

# Opinion Dynamics for Agents with Resource Limitations

Pavan Tallapragada

Associate Professor

Robert Bosch Centre for Cyber Physical Systems

Indian Institute of Science



Decisions, Games, and Evolution  
ICTS

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# Opinion Dynamics



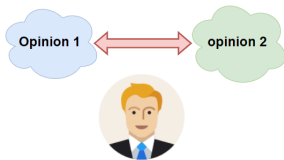
Consensus



Clustering



Polarization



Oscillations



Social Power



Opinion manipulation

# Opinion and Influence Modeling

Opinion of an agent on a single topic is a scalar  $x \in \mathbb{R}$

$$x < 0$$

Opposing  
Opinion

$$x = 0$$

Neutral  
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$$x > 0$$

Favoring  
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*Magnitude denotes the extent to which the agent favors or opposes the topic*

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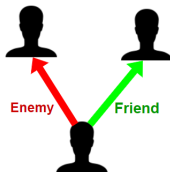
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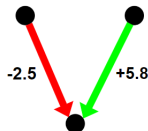
Favoring  
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*Magnitude denotes the extent to which the agent favors or opposes the topic*

- Modeling of social connections among the agents using a graph.*



Social ties



Directed signed graph

# Literature Review

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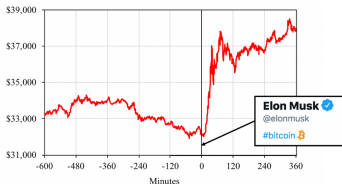
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- **Utility Maximization and Game theoretic approach:** Ghaderi and Srikant (2014), Etesami et.al (2019), Etesami and Basar (2015), Etesami (2022), Park et.al (2022).

# Motivation for Our Work

- Resources influence the opinions and social power.

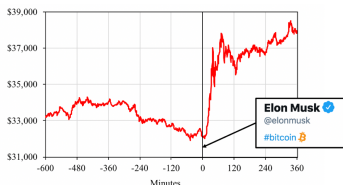
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- Resources restrict extreme opinions.

## House Auction



**Bidder 1**

**Bid: \$150k**

**Budget: \$300k**



**Bidder 2**

**Bid: \$180k**

**Budget: \$500k**



**Bidder 3**

**Bid: \$200k**

**Budget: \$1M**

# Problem Formulation

- $\mathcal{A} := \{1, \dots, n\}$  : Set of agents discussing on a single topic.
- $z_i(t) \in \mathbb{R}$  : Opinion of  $i \in \mathcal{A}$  at time  $t$ .
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$$\dot{\mathbf{z}}_i := S_i(z_i) + C_i(\mathbf{z}), \forall i \in \mathcal{A}, \quad (OD)$$

where

$$S_i(z_i) := -w_i [z_i - p_i] - \frac{z_i^3}{r_i} \quad (Self),$$

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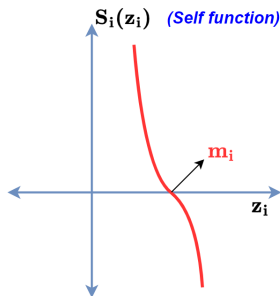
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$S_i(\cdot)$  is strictly  $\downarrow$  and  
 $S_i(\mathbf{m}_i) = 0$

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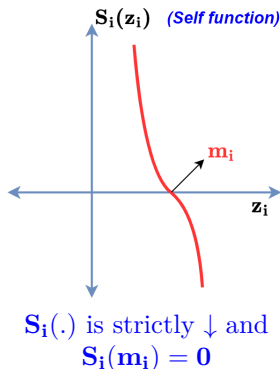
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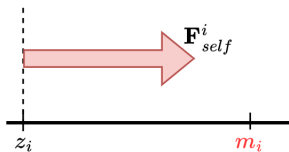


- $\mathcal{E}$  : Set of equilibria of OD.
- $\mathbf{m}_i \in \mathbb{R}$  is the opinion agent  $i$  would hold, if  $i$  was “isolated” from social influence.

# Force analogy

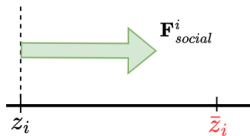
- Force due to  $S_i(z_i)$  :

Attractive force towards  $m_i$

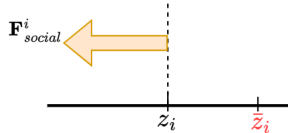


- Force due to  $C_i(\bar{z}_i)$  :

$c_i > 0$ : Attractive force towards  $\bar{z}_i$



$c_i < 0$ : Repulsive force away from  $\bar{z}_i$



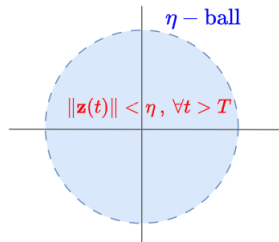


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## Theorem (Ultimate Boundedness)

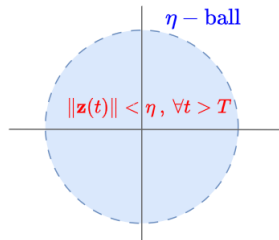
Under OD, the opinions are **ultimately bounded**.



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## Theorem (Existence and Uniqueness of Equilibrium)

Let  $\mathcal{N}_i^e$  denote the **set of enemies** of  $i$ . Suppose

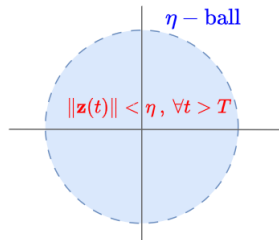
$$w_i > \sum_{k \in \mathcal{N}_i^e} 2|a_{ik}|; \forall i \in \mathcal{A} \quad (\text{WAR: Weak antagonistic relations})$$

Then,

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Then, there exists **an unique GES equilibrium point**  $\mathbf{z}^* \in \mathcal{E}$ .

# Ultimate Bound (No Antagonistic Relations)

- Let  $\mathcal{M} := [m_{\min}, m_{\max}]$ . Also define  $\mathcal{V}_{\max} := \{i \in \mathcal{A} \mid m_i = m_{\max}\}$  and  $\mathcal{V}_{\min} := \{i \in \mathcal{A} \mid m_i = m_{\min}\}$ .

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- $\mathbf{z}(t)$  converges to **positive invariant** set  $\mathcal{M}^n$ .
- $\mathbf{z}(t)$  converges to **interior** of  $\mathcal{M}^n$  in **finite time** iff every agent in  $\mathcal{V}_{\max}$  and  $\mathcal{V}_{\min}$  is directly/indirectly influenced by at least one agent in  $\mathcal{A} \setminus \mathcal{V}_{\max}$  and  $\mathcal{A} \setminus \mathcal{V}_{\min}$  resp.
  - *This condition is satisfied if the social network is strongly connected.*



# Consensus

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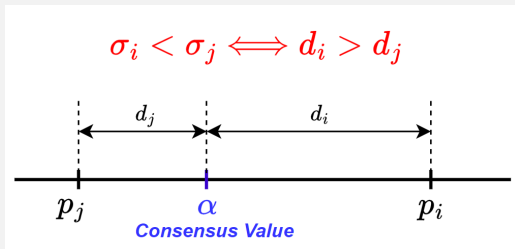
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## Lemma (Dominance at Consensus)

For each  $i \in \mathcal{A}$ , define (dominance weight)  $\sigma_i := w_i r_i$ . Then for any  $i, j \in \mathcal{A}$ ,



**Opinion Game:**  $\mathcal{G} = \langle \mathcal{A}, (\mathbb{R})_{i \in \mathcal{A}}, (U_i)_{i \in \mathcal{A}} \rangle$

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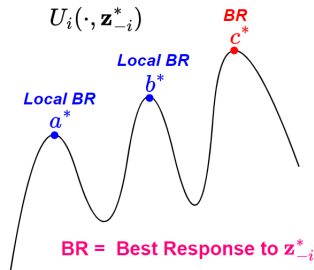
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$$U_i(z_i^*, \mathbf{z}_{-i}^*) \geq U_i(z_i, \mathbf{z}_{-i}^*), \quad \forall z_i \text{ s.t. } |z_i^* - z_i| \leq \rho_i$$

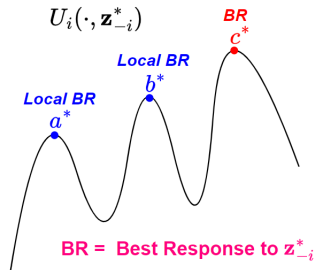
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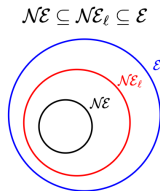
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# Nash Equilibria

## Theorem

Suppose  $\mathbf{z}^* = (z_i^*, \mathbf{z}_{-i}^*) \in \mathcal{E}$ . Let  $\eta$  denote the **ultimate bound** on opinions. For each  $i \in \mathcal{A}$ , define

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- If the above inequality is strict, then  $\mathbf{z}^* \in \mathcal{NE}_\ell$ .

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$$\tau_i := \frac{-r_i}{3} \left[ w_i + \sum_{k \in \mathcal{A}/\{i\}} a_{ik} \right] =: \frac{-r_i d_i^{in}}{3}. \text{ Then,}$$

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- *Opinions under OD converge to this unique Nash Eq.*

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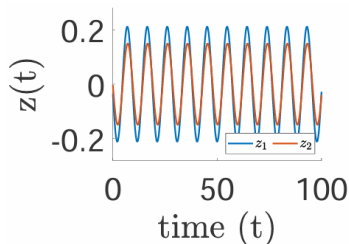
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- *Everyone is satisfied with the outcome!*

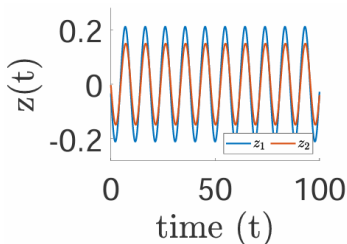


# Periodic Behavior of Opinions: Two Agent Dynamics



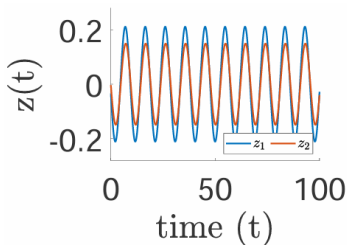
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- The model could exhibit oscillations in presence of contrarians.
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- A sufficient condition for Hopf Bifurcation to exist.

# Summary

- OD derived out of the utility
- Limited resources prevent extreme opinions
- Game theoretic analysis
- Ongoing work: Multi-topic OD under hard resource constraints

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