

Module – 2.4

OPTIMIZATION IN FINITE DIMENSIONAL VECTOR SPACES: AN OVERVIEW

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A CLASSIFICATION – MINIMUM AND MAXIMUM

- Let $f: \mathbb{R}^n \rightarrow \mathbb{R}$ and $f \in C^2(\mathbb{R}^n)$
- f is twice continuously differentiable functional

$$\min_{x \in \mathbb{R}^n} f(x) = \max_{x \in \mathbb{R}^n} \{-f(x)\}$$

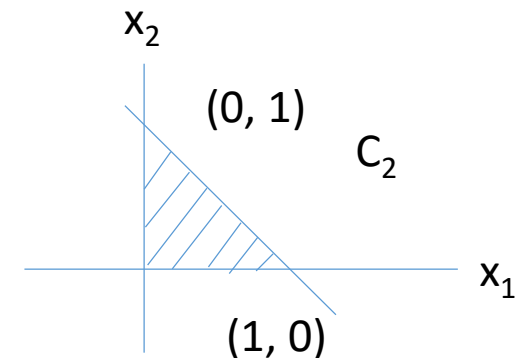
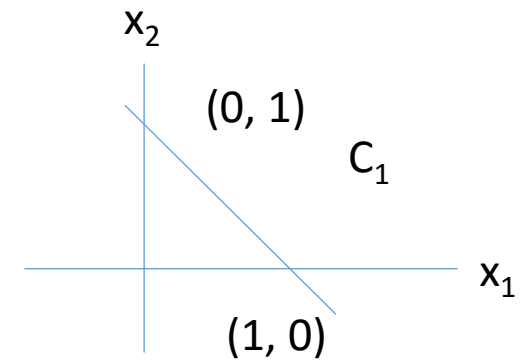
- Consider only minimization

A CLASSIFICATION – UNI VS MULTI MODAL

- $N_\varepsilon(x) = \{ y \in R^n \mid \|y - x\| \leq \varepsilon \} \subset R^n$ called ε – neighborhood
- If $x^* \in R^n$ is such that $f(x^*) \leq f(y)$ for all $y \in N_\varepsilon(x^*)$, then x^* is a local minimum
- If $x^* \in R^n$ is such that $f(x^*) \leq f(y)$ for all y , then x^* is a global minimum
- A function that has a unique minimum is a unimodal function
- Otherwise, it is a multimodal function
- $f(x) = x(x^2 - 1)$ is multimodal, $f(x) = x^2$ is unimodal

A CLASSIFICATION – CONSTRAINED VS UNCONSTRAINED

- Let $C \subset \mathbb{R}^n$ defined by a set of equations or inequalities
- $C_1 = \{ x \in \mathbb{R}^2 \mid x_1 + x_2 = 1 \}$
- $C_2 = \{ x \in \mathbb{R}^2 \mid x_1 \geq 0, x_2 \geq 0, x_1 + x_2 \leq 1 \}$
- Let $f(x) = x_1^2 + x_2^2$
- $\text{Min}_{x \in \mathbb{R}^2} f(x)$ is an unconstrained minimization
- $\text{Min}_{x \in C_1} f(x)$ is a constrained minimization with equality constraints
- $\text{Min}_{x \in C_2} f(x)$ is a constrained minimization with inequality constraints
- Linear and non linear programming deal with minimization under inequality constraints
- In this course we will deal with unconstrained and equality constrained minimization problem only



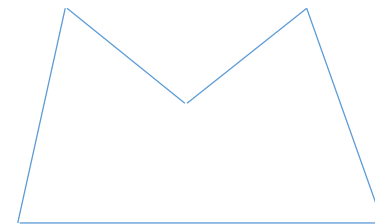
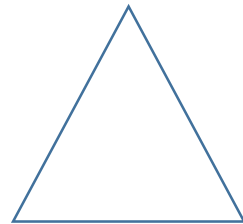
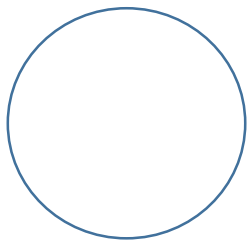
A CLASSIFICATION – UNI VS MULTI-OBJECTIVE OPTIMIZATION

- If $f: \mathbb{R}^n \rightarrow \mathbb{R}$ is the only function to be minimized, it is known as uni-objective minimization
- If $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$ with $f(x) = (f_1(x), f_2(x), \dots, f_m(x))^T$ where we want to minimize some and maximize others, it is called multi-objective optimization
- Automobile design – Maximize fuel efficiency, minimize cost, maximize safety and comfort is an example of multi-objective optimization
- In this course we only deal with uni-objective minimization

ROLE OF CONVEXITY IN MINIMIZATION

- Let S be a subset of \mathbb{R}^n
- S is called a convex set if for every pair of points x and y in S , the points along the line segment joining x and y are also in S

$$\alpha x + (1 - \alpha)y \in S \text{ if } x, y \in S$$

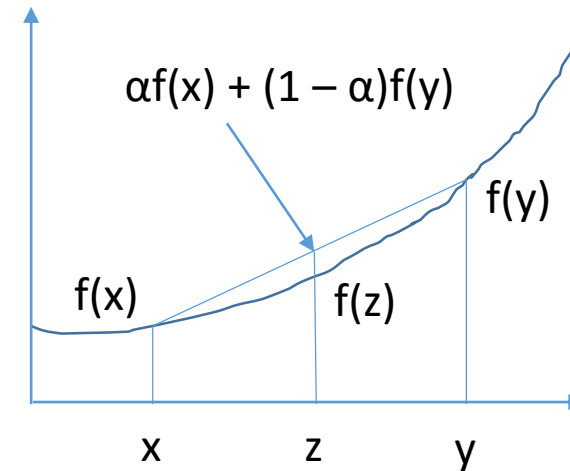


Convex sets

Non-Convex set

ROLE OF CONVEXITY IN MINIMIZATION

- Let S be a convex set in \mathbb{R}^n and let $x, y \in S$
- A function $f: S \rightarrow \mathbb{R}$ is said to be a convex function if
$$f(\alpha x + (1 - \alpha)y) \leq \alpha f(x) + (1 - \alpha)f(y)$$
for all $\alpha \in [0, 1]$
- A convex function lies below the chord
- Let $z = \alpha x + (1 - \alpha)y$
- If $f(x)$ is convex, $-f(x)$ is concave
- $f(x) = x^2$ is convex but $g(x) = x^3$ is not



CHARACTERIZATION OF CONVEXITY

- If $f \in C^1(S)$ be a continuously differential function defined on a convex set S
- f is convex if and only if, for $x, y \in S$
$$f(y) \geq f(x) + (y - x)^T \nabla_x f(x) \rightarrow \text{curve lies above the tangent}$$
- $f(x)$ is strictly convex if strict inequality holds
- If $f \in C^2(S)$ be twice continuously differentiable function on a convex set S
- Then f is convex if and only if the Hessian $\nabla_x^2 f(x)$ is positive semi-definite. f is strictly convex if $\nabla_x^2 f(x)$ is positive definite

CONVEXITY AND UNIMODALITY

- $f: S \rightarrow \mathbb{R}$ and S is a convex set
- Then f has a unique minimum
- If $f \in C^2(S)$, then at this minimum $\nabla_x f(x) = 0$ and $\nabla_x^2 f(x)$ is positive definite
- $f = x^T A x - b^T x$ is a typical convex function in $C^2(\mathbb{R}^n)$ when A is symmetric and positive definite

CONDITIONS FOR UNCONSTRAINED MINIMUM

- $f: \mathbb{R}^n \rightarrow \mathbb{R}$ and $f \in C^2(\mathbb{R})$
- A necessary condition for the minimum is that at the minimum
$$\nabla_x f(x) = 0 \rightarrow \text{Gradient Vanishes}$$
- A sufficient condition for the minimum is that at the minimum
$$\nabla_x^2 f(x) \text{ SPD} \rightarrow \text{Hessian is a (symmetric) positive definite matrix}$$

EQUALITY CONSTRAINED MINIMUM - ELIMINATION

- Method of elimination : Illustration
- Maximize $A = ab$ when $2(a + b) = L$ is fixed
- Eliminate b in A : $b = \frac{L}{2} - a$ and $A = a(\frac{L}{2} - a) = \frac{La}{2} - a^2$
$$\frac{dA}{da} = \frac{L}{2} - 2a \text{ and } \frac{d^2A}{da^2} = -2 < 0$$
- At the maximum $a^* = \frac{L}{4}$ and $b^* = \frac{L}{4}$ and $A_{\max} = \frac{L^2}{16}$

EQUALITY CONSTRAINED MINIMIZATION – LAGRANGIAN MULTIPLIER

- Method of Lagrangian multiplier

- Let $g: \mathbb{R}^n \rightarrow \mathbb{R}^m$ and $g \in C^2(\mathbb{R}^n)$

- $g(x) = (g_1(x), g_2(x), \dots, g_m(x))^T$

- Min $f(x)$ when $g(x) = b$ where $b \in \mathbb{R}^m$

- Define the Lagrangian

$$L(x, \lambda) = f(x) + \lambda^T(b - g(x))$$

- $\lambda \in \mathbb{R}^m$ is the vector of undetermined Lagrangian multiplier

- At the minimum:

$$\nabla_x L(x, \lambda) = \nabla_x f(x) - \sum_{i=1}^m \lambda_i \nabla_x g_i(x) = 0$$

$$\nabla_x L(x, \lambda) = b - g(x) = 0$$

- A necessary condition for the minimum is that at the minimum the gradient $\nabla_x f(x)$ must be a linear combination of the gradients of the constraints

SUFFICIENT CONDITION FOR EQUALITY CONSTRAINTS

- The Hessian of $L(x, \lambda)$ is given by

$$\nabla_x^2 L(x, \lambda) = \nabla_x^2 f(x) - \sum_{i=1}^m \lambda_i \nabla_x^2 g_i(x)$$

- Let $T = \{ y \in \mathbb{R}^n \mid y^\top \nabla g_i(x) = 0, 1 \leq i \leq m \}$
- T consists of all vectors that are orthogonal to $\nabla g_i(x), 1 \leq i \leq m$.
Indeed, T is the tangent plane to $g_i(x), 1 \leq i \leq m$
- Let x^* be such that there exists $\lambda^* \in \mathbb{R}^m$ with
 - a) $\nabla_x f(x) = \sum_{i=1}^m \lambda_i^* \nabla_x g_i(x)$
 - b) $\nabla_x^2 L(x^*, \lambda)$ is positive definite on T
 $y^\top \nabla_x^2 L(x^*, \lambda) y > 0$ for all $y \in T$Then, x^* is a relative constrained minimum

ILLUSTRATION – EQUALITY CONSTRAINT

- Let $n = 2$, $f(x) = x_1 + x_1x_2 + 3x_2^2$ - to be minimized

$$g(x) = x_1 + 2x_2 - 3 = 0 - \text{constraint}$$

- $L(x, \lambda) = (x_1 + x_1x_2 + 3x_2^2) - \lambda(x_1 + 2x_2 - 3)$

- First-order necessary condition:

$$\nabla_x f(x) - \lambda \nabla_x g(x) = \begin{bmatrix} 1 + x_2 - \lambda \\ x_1 + 6x_2 - 2\lambda \end{bmatrix} = 0$$

$$x_1 + 2x_2 - 3 = 0$$

- Solution: $x_1^* = 4, x_2^* = -\frac{1}{2}, \lambda^* = \frac{1}{2}$

- $\nabla_x^2 f(x) = \begin{bmatrix} 0 & 1 \\ 1 & 6 \end{bmatrix}, \nabla_x^2 g(x) = \begin{bmatrix} 0 & 0 \\ 0 & 0 \end{bmatrix}, \nabla_x^2 L(x, \lambda) = \nabla_x^2 f(x)$

- $\nabla_x f(x^*) = \frac{1}{2} \begin{pmatrix} 1 \\ 2 \end{pmatrix}, \Rightarrow T = \left\{ \frac{2\alpha}{\sqrt{5}} \begin{pmatrix} -2 \\ 1 \end{pmatrix} \mid \alpha \in \mathbb{R} \right\}$

- $\frac{2\alpha}{\sqrt{5}}(-2, 1) \begin{bmatrix} 0 & 1 \\ 1 & 6 \end{bmatrix} \frac{2\alpha}{\sqrt{5}} \begin{pmatrix} -2 \\ 1 \end{pmatrix} = \frac{8\alpha^2}{5} > 0$

- Hence $(x_1^* = 4, x_2^* = -\frac{1}{2})$ is a constrained minimum

PENALTY FUNCTION METHOD – EQUALITY CONSTRAINT

- Let $f: \mathbb{R}^n \rightarrow \mathbb{R}$, $g: \mathbb{R}^n \rightarrow \mathbb{R}^m$
- Minimize $f(x)$ when $g(x) = b$
- Consider $P_\alpha(x) = f(x) + \alpha g^T(x)g(x) = f(x) + \frac{\alpha}{2} \sum_{i=1}^m g_i^2(x)$
- $\alpha > 0$ is called the penalty constant
- $\nabla_x P_\alpha(x) = \nabla_x f(x) + \alpha D_x^T(g)g(x)$
where $D_x(g) \in \mathbb{R}^{m \times n}$ is the Jacobian of $g(x)$
- Let $x^*(\alpha)$ be the solution of $\nabla_x P_\alpha(x) = 0$

PENALTY FUNCTION METHOD – EQUALITY CONSTRAINT

- It can be shown $\lim_{\alpha \rightarrow \infty} x^*(\alpha) \rightarrow x^*$, the constrained minimum
- Rewrite

$$\begin{aligned}\nabla_x P_\alpha(x) &= \nabla_x f(x) + \sum_{i=1}^m \nabla_x g_i(x) [\alpha g_i(x)] \\ &= \nabla_x f(x) + \sum_{i=1}^m \nabla_x g_i(x) \lambda_i(\alpha)\end{aligned}$$

where $\lambda_i(\alpha) = \alpha g_i(x)$, $1 \leq i \leq m$ plays the role of the Lagrangian multiplier

- It can be shown $\lim_{\alpha \rightarrow \infty} \lambda_i(\alpha) = \lambda_i^*$, the value of the Lagrangian multiplier at the minimum

ILLUSTRATION

- $n = 2$, $f(x) = x_1^2 + x_2^2$ - to be minimized

$$g(x) = x_1 + x_2 - 1 \text{ - constraint}$$

- $P_\alpha(x) = x_1^2 + x_2^2 + \frac{\alpha}{2} [x_1 + x_2 - 1]^2$

$$\nabla_x P_\alpha(x) = \begin{bmatrix} x_1(2 + \alpha) + \alpha x_2 - \alpha \\ \alpha x_1 + (2 + \alpha) - \alpha \end{bmatrix} = 0$$

$$\Rightarrow x^*(\alpha) = \left(\frac{1}{2 + \alpha^{-1}}, \frac{1}{2 + \alpha^{-1}} \right)^\top$$

- Multiplier $\lambda(\alpha) = \alpha g(x^*(\alpha)) = \frac{1}{1 + \alpha^{-1}}$

- As $\alpha \rightarrow \infty$, $x^* = \left(\frac{1}{2}, \frac{1}{2} \right)^\top$ and $\lambda^* = 1$ which is the minimizer obtained using Lagrangian multiplier method

STRONG VS WEAK CONSTRAINED FORMULATION

- Min $f(x)$ when $g(x) = b$
- Lagrangian multiplier method is called strong constraint formulation which forces the exact equality condition
- Penalty function method is called weak constraint formulation which only forces approximate equality depending on the value of α , the solution is more closer the constraint and $\alpha \rightarrow \infty$, the constraint is exactly satisfied
- We will use both of these formulations

EXERCISES

5.1) Plot $f(x) = x(x^2 - 1)$ for $-2 \leq x \leq 2$ and identify the minima and maxima

5.2) Let $f(x)$ has a minimum at x^* then show that $af(x)$, $f(x) + c$, and $af(x) + c$ all have a minimum at x^*

5.3) Find the minimizer of $f(x) = x_1^2 + x_2^2$ when $x_1 + x_2 = 1$ using Lagrangian multiplier method

5.4) Find the x that minimizes $\frac{\alpha}{2} \|x\|^2$ under the constraint $Z = Hx$ using (a) Lagrangian multiplier and (b) Penalty function method

5.5) Find the minimizer of

1) $f_1(x) = (Z - Hx)^T W (Z - Hx)$

2) $f(x) = f_1(x) + (x - x_b)^T B^{-1} (x - x_b)$

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- D. G. Luenberger (1969) Otimization in Vector Spaces, *Wiley*
- D. G. Luenberger (1973) Introduction to Linear and Nonlinear Programming, *Addison Wesley*
- S. G. Nash and A. Sofer (1996) Linear and Nonlinear Programming, *McGraw Hill*