

Decisions, Games, and Evolution Bangalore 2025

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Example 1: Mating behavior among male lizards (Sinervo & Lively 1996)

• Among side-blotched lizards (*Uta Stansburiana*), there are three male morphs. One can distinguish them by their throat color: yellow, blue, orange.

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B. Sinervo & C. M. Lively



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Question: How can we make sense of this coexistence of different morphs?

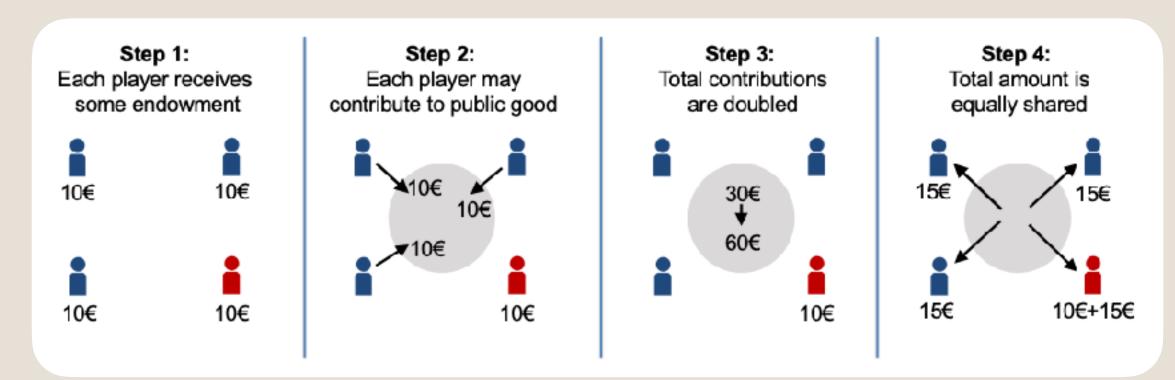
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Example 2: Cooperation and punishment among humans (Fehr & Gächter 2000)

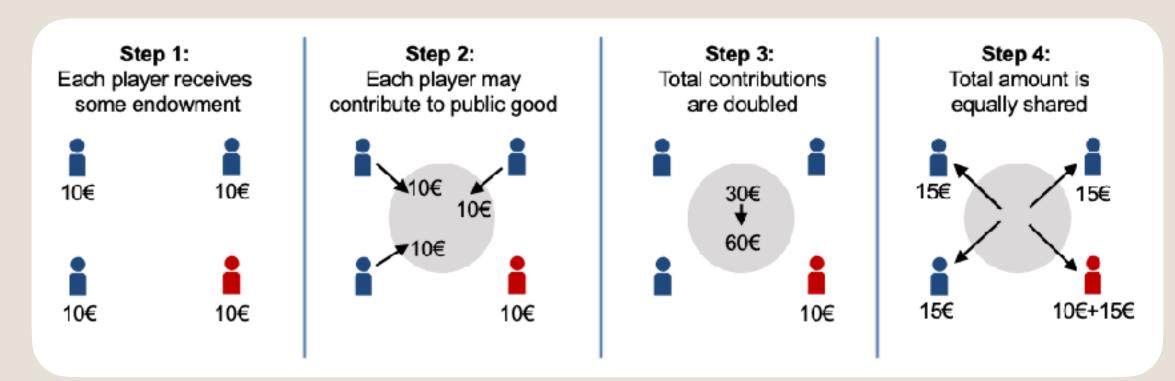
Cooperation and Punishment in Public Goods Experiments



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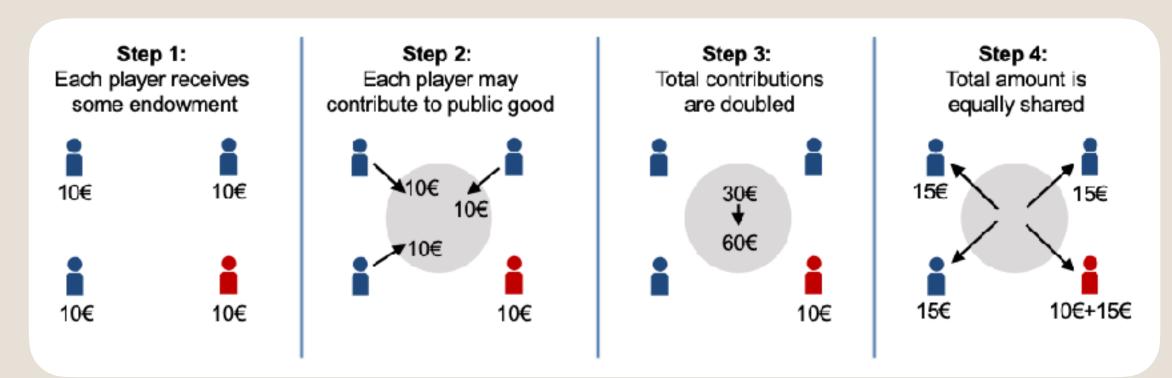
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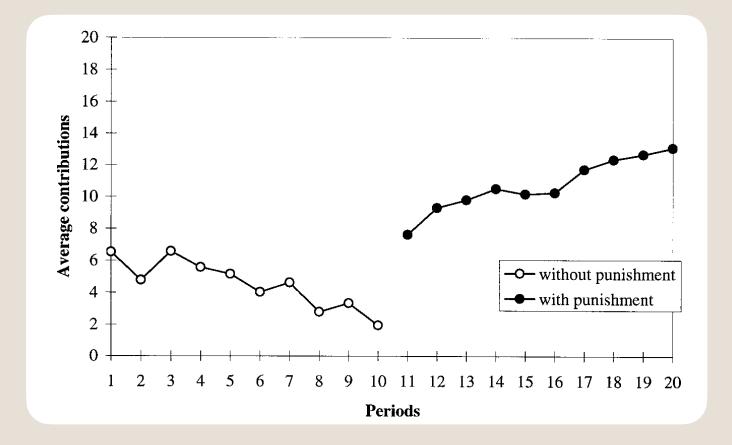
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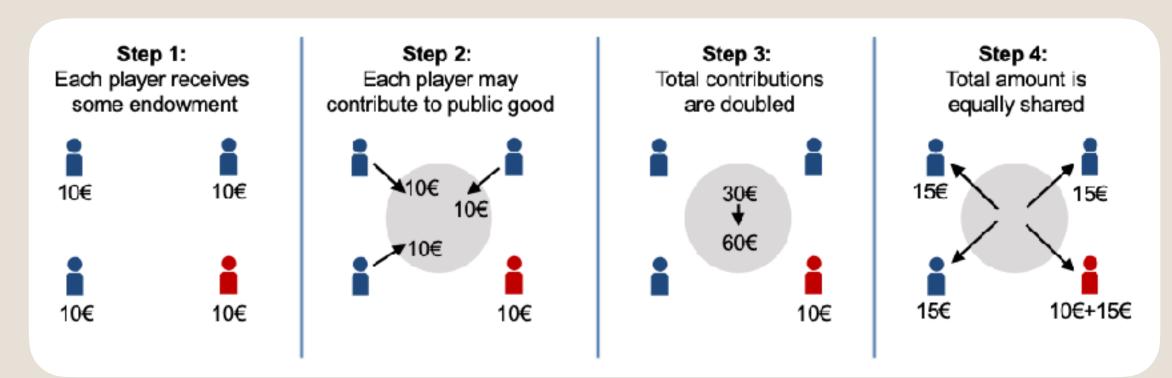


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 - Without punishment, cooperation goes down
 - With punishment, cooperation goes up even if participants never meet again (at least in most Western countries)

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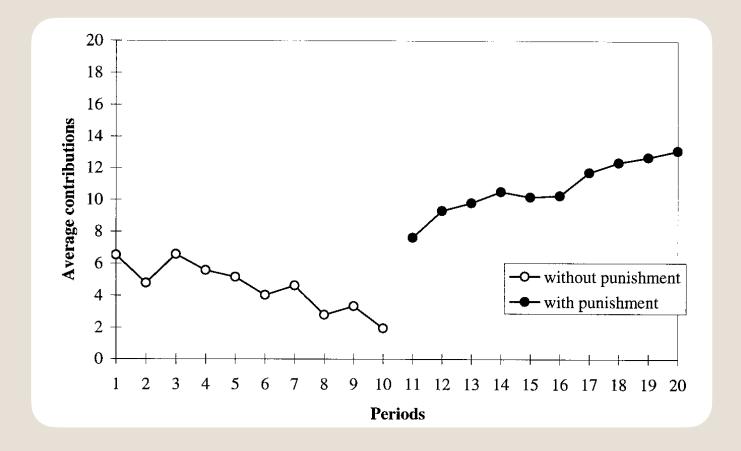


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Question: How can we make sense of these behaviors?

Cooperation and Punishment in Public Goods Experiments



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JOB MARKET SIGNALING *

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JOB MARKET SIGNALING * Michael Spence

Mate Selection—A Selection for a Handicap

AMOTZ ZAHAVI



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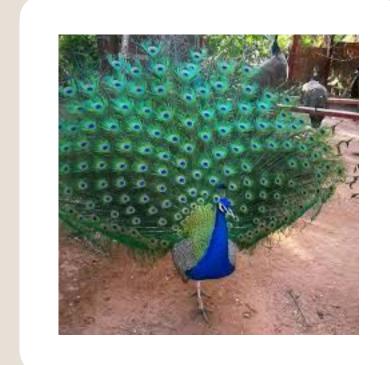
- Individuals sometimes invest in something without getting a direct return
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Question: It has been suggested that these investments can be worthwhile when they act as (costly) signals. But how exactly do such signals work?

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Interesting observation

Not in all examples the respective behaviors and traits are consciously chosen.

An overview

Today's class (March 11, 2025)

 An introduction to evolutionary game theory (Replicator dynamics, games in finite populations)

Tomorrow's classes (March 12, 2025)

- Evolution of cooperation & direct reciprocity
- Social norms & indirect reciprocity

Thursday's class (March 13, 2025)

• Some current research: Reciprocity in complex environments

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Such games can be represented by a (bi)-matrix

Action 1 ... Action
$$n$$

Action 1 a_{11}, b_{11} ... a_{1n}, b_{1n}

... Action m a_{m1}, b_{m1} ... a_{mn}, b_{mn}

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Definition: Dominated strategies

A pure strategy \mathbf{e}_i for player 1 is called (strictly) dominated if there is a (possibly mixed) strategy \mathbf{x} for player 1 that yields a better payoff, irrespective of the co-player's strategy \mathbf{e}_j , $\pi_1(\mathbf{e}_i,\mathbf{e}_i) < \pi_1(\mathbf{x},\mathbf{e}_i)$ for all \mathbf{e}_j .

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Definition: Nash equilibrium

A strategy profile (x^*,y^*) is called a Nash equilibrium if the following two conditions hold:

$$\pi_1(\mathbf{x}, \mathbf{y}^*) \le \pi_1(\mathbf{x}^*, \mathbf{y}^*) \text{ for all } \mathbf{x} \in S_m.$$
 (1.13.1)

$$\pi_2(\mathbf{x}^*, \mathbf{y}) \le \pi_2(\mathbf{x}^*, \mathbf{y}^*) \text{ for all } \mathbf{y} \in S_n.$$
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Dominance solvability and the Nash equilibrium concept appear to make strong assumptions on cognitive abilities. In the following, we explore an approach to game theory that avoids these assumptions.

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$$f_H = \frac{b-c}{2}x + b(1-x) \quad \text{and} \quad f_D = 0 \cdot x + \frac{b}{2}(1-x)$$

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• The larger the cost of serious injuries, the more doves we would expect.

Example 1.2: Hawk-Dove as a classical game

• We could have also interpreted this interaction as a classical game with payoff matrix

	Hawk	Dove
Hawk	(b-c)/2, (b-c)/2	b , 0
Dove	0, b	b/2, b/2



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• Consider an infinitely large population

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Remark 1.3: Introducing matrix games for populations

- Consider an infinitely large population
- Individuals in that population can have one of n different traits ("strategies"). Let $\mathbf{x} = (x_1, \dots, x_n)^T$ describe the trait distribution in the population.

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• Similarly, the population's average fitness is

$$\bar{f} = \sum_{i=1}^{n} x_i f_i = \mathbf{x}^T A \mathbf{x}$$

Definition 1.4: Replicator equation / Replicator dynamics

The replicator equation is the system of ordinary differential equations

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$$\sum_{i=1}^{n} x_i(0) = 1 \quad \Rightarrow \quad \sum_{i=1}^{n} x_i(t) = 1 \quad \forall t$$

2. All boundary faces of S_n are also invariant:

$$x_i(0) = 0 \Rightarrow x_i(t) = 0 \ \forall t$$

3. Dominated traits go extinct:

$$f_i(\mathbf{x}) < f_j(\mathbf{x}) \ \forall \mathbf{x} \ \text{and} \ \mathbf{x}(0) \in \text{int}(S_n) \ \Rightarrow \ \lim_{t \to \infty} x_i(t) = 0$$

Definition 1.4: Replicator equation / Replicator dynamics

The replicator equation is the system of ordinary differential equations

$$\dot{x}_i = x_i (f_i(\mathbf{x}) - \bar{f}(\mathbf{x})).$$

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Proof sketch.

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$$\left(\frac{\dot{x}_i}{x_j}\right) = \frac{\dot{x}_i x_j - x_i \dot{x}_j}{x_i^2}$$

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$$\begin{pmatrix} \dot{x}_i \\ x_j \end{pmatrix} = \frac{\dot{x}_i x_j - x_i \dot{x}_j}{x_j^2}$$

$$= \frac{x_i x_j (f_i(\mathbf{x}) - \bar{f}(\mathbf{x})) - x_i x_j (f_j(\mathbf{x}) - \bar{f}(\mathbf{x}))}{x_j^2}$$

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= \frac{x_i x_j (f_i(\mathbf{x}) - \bar{f}(\mathbf{x})) - x_i x_j (f_j(\mathbf{x}) - \bar{f}(\mathbf{x}))}{x_j^2}$$

$$= \frac{x_i}{x_j} \left(f_i(\mathbf{x}) - f_j(\mathbf{x}) \right) < -\delta \left(\frac{x_i}{x_j} \right)$$

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3. Consider the fraction x_i/x_i .

$$\begin{pmatrix} \dot{x}_i \\ x_j \end{pmatrix} = \frac{\dot{x}_i x_j - x_i \dot{x}_j}{x_j^2} \\
= \frac{x_i x_j (f_i(\mathbf{x}) - \bar{f}(\mathbf{x})) - x_i x_j (f_j(\mathbf{x}) - \bar{f}(\mathbf{x}))}{x_j^2}$$

$$= \frac{x_i}{x_j} \left(f_i(\mathbf{x}) - f_j(\mathbf{x}) \right) < -\delta \left(\frac{x_i}{x_j} \right)$$

Therefore, the fraction x_i/x_j decreases exponentially.

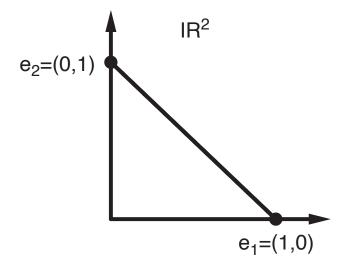
Remark 1.6: On representing the unit simplex

$$S_n = \left\{ \mathbf{z} \in \mathbb{R}^n : z_i \ge 0 \text{ for all } i \text{ and } \sum_{i=1}^n z_i = 1 \right\}$$

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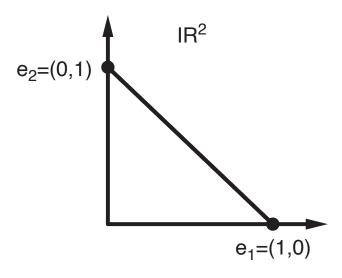
$$n = 2$$

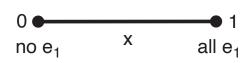


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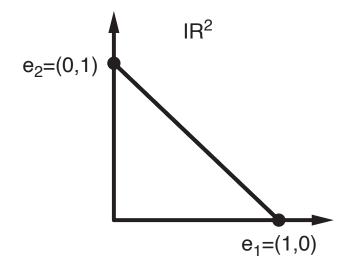




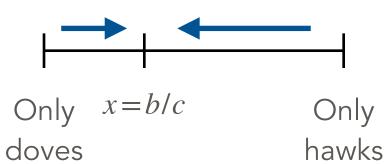
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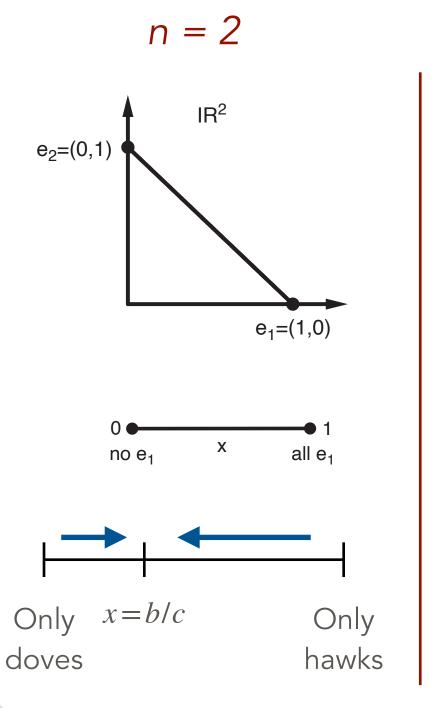


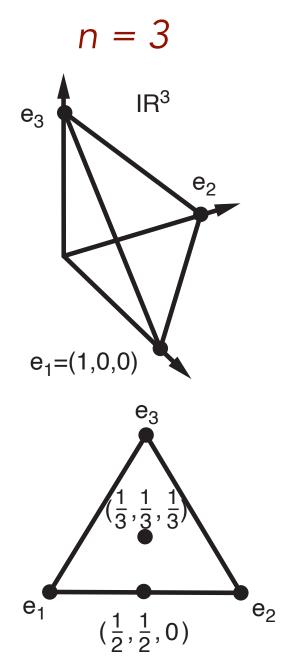




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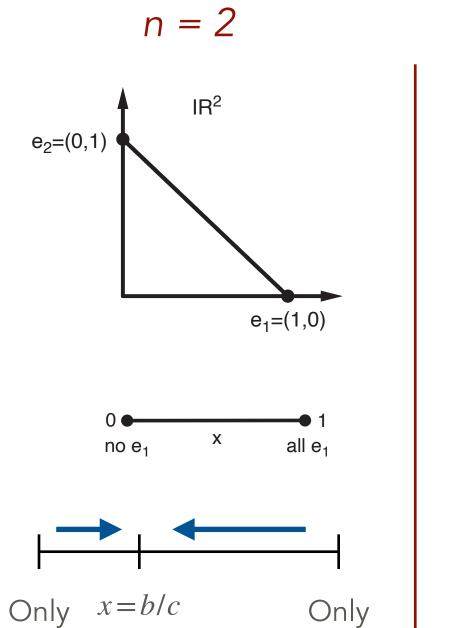




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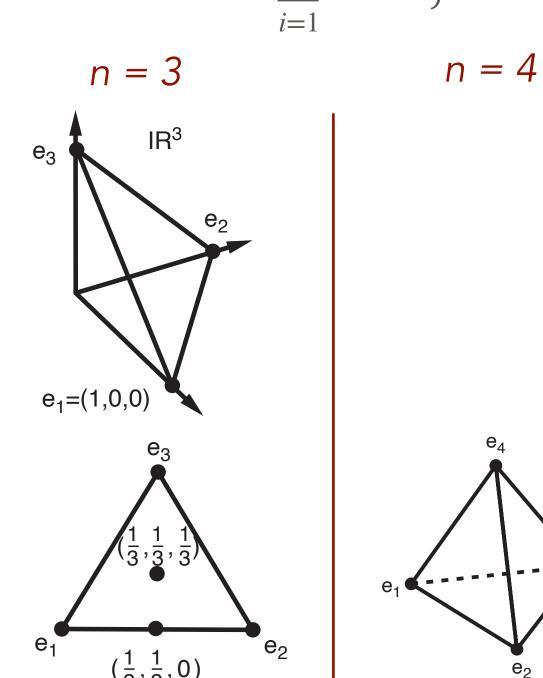
Consider the replicator equation $\dot{x}_i = x_i (f_i(\mathbf{x}) - \bar{f}(\mathbf{x}))$. For a game with n strategies in total, this is, in principle, an *n*-dimensional system. However, we are only interested in those orbits on the unit simplex:

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doves

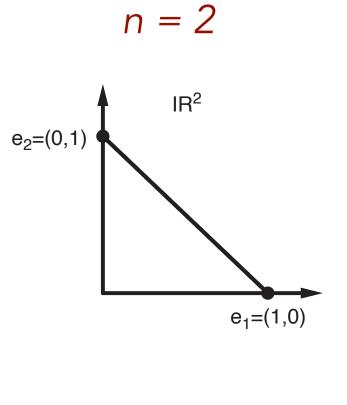
hawks

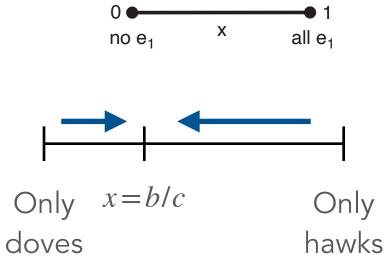


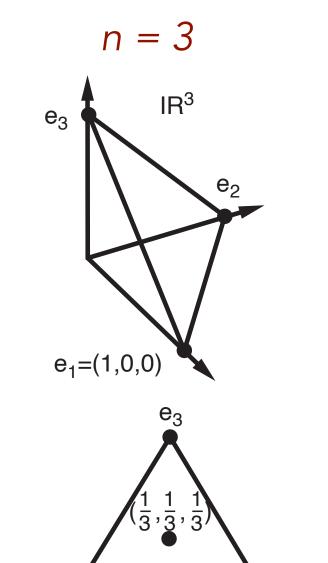
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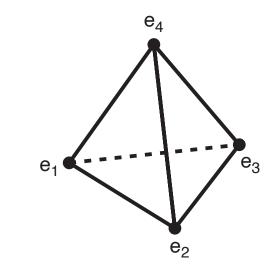
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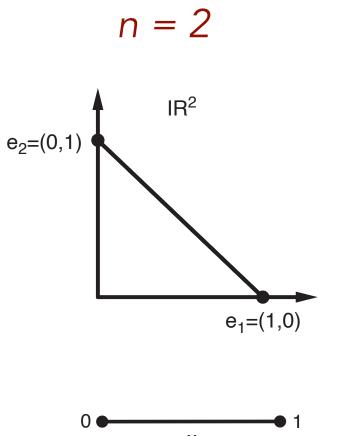
Remark 1.7: A classification of 2x2 games

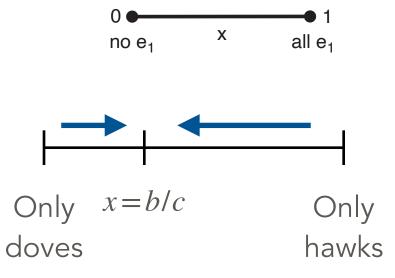
To get some intuition, let us analyze the simplest non-trivial case: a symmetric game with two strategies:

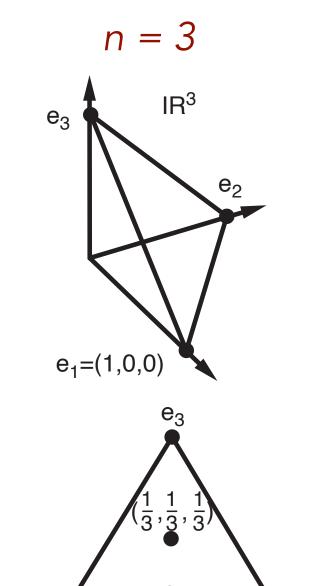
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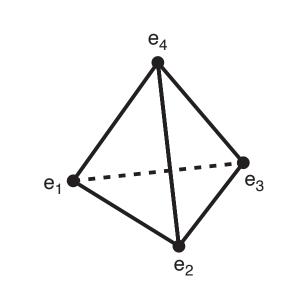
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n = 4

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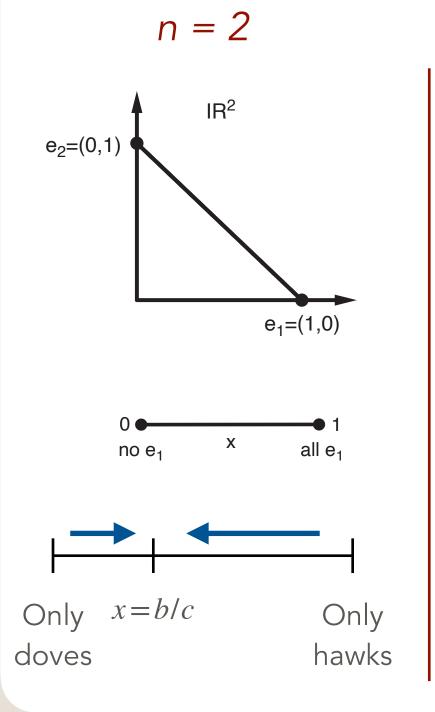
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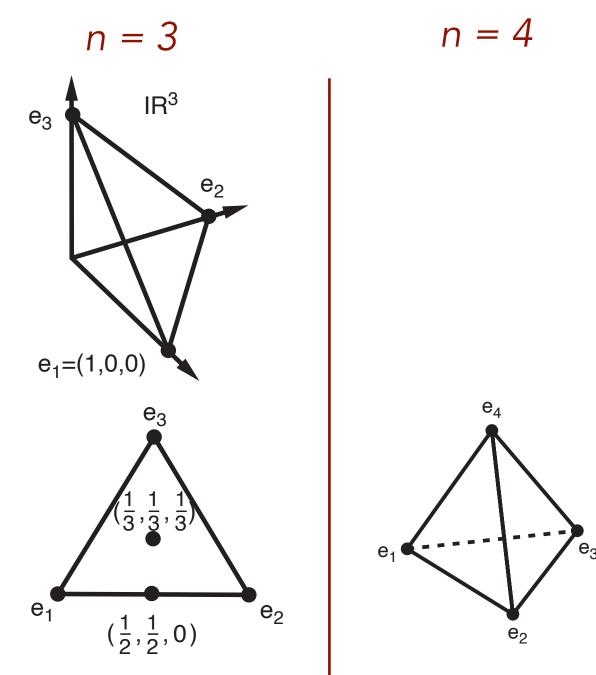
	Action 1	Action 2
Action 1	а	b
Action 2	C	d

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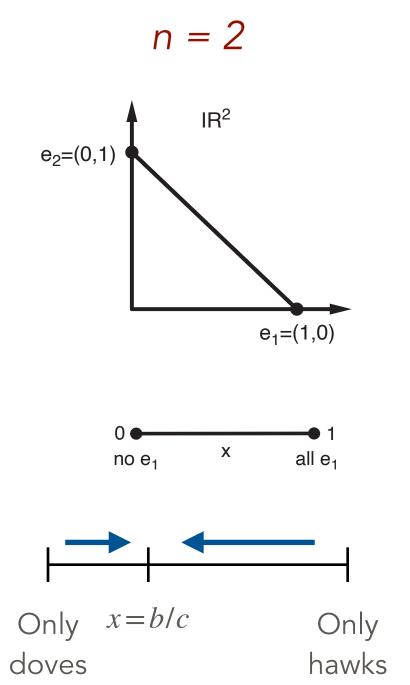
	Action I	Action 2
Action 1	а	b
Action 2	C	d

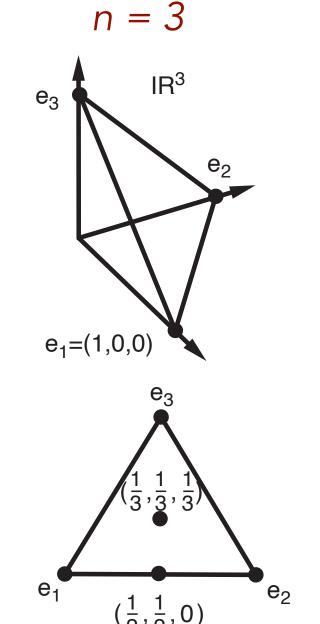
We can represent the replicator equation as a 1-dim. system. Let x be the proportion of individuals who use Action 1, and 1-x is the proportion of individuals who use Action 2.

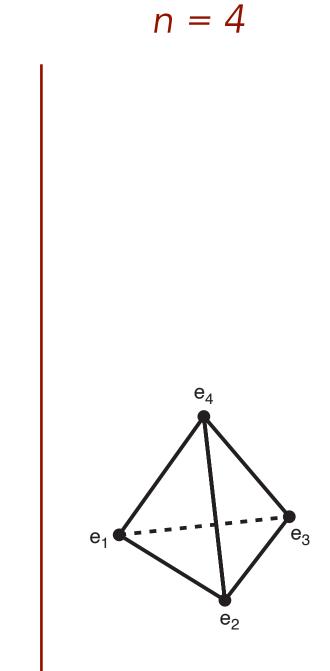
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	Action 1	Action 2
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Action 2	С	d

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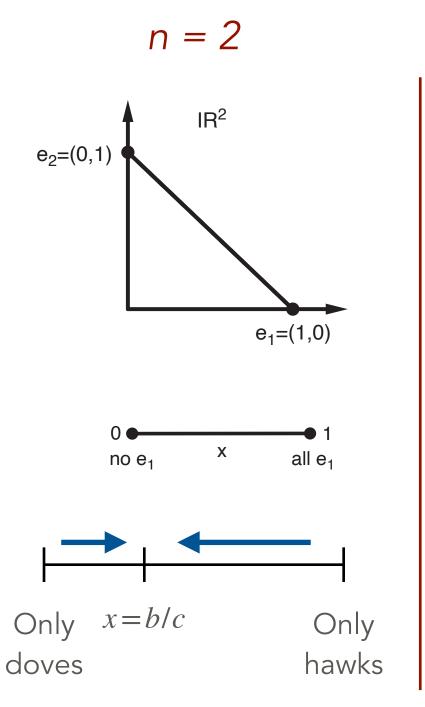
The fitnesses are

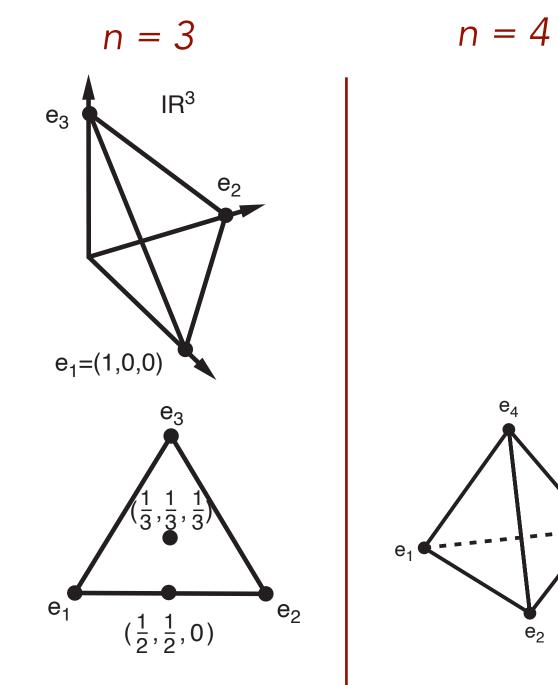
$$f_1(x) = ax + b(1-x)$$
 and $f_2(x) = cx + d(1-x)$

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The fitnesses are

$$f_1(x) = ax + b(1 - x)$$
 and $f_2(x) = cx + d(1 - x)$

Replicator dynamics takes the form:

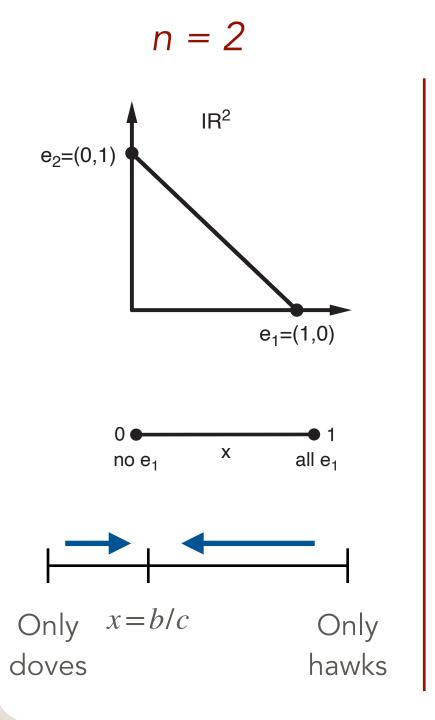
$$\dot{x} = x \left(f_1(x) - \bar{f}(x) \right) = x \left(f_1(x) - x f_1(x) - (1 - x) f_2(x) \right)$$

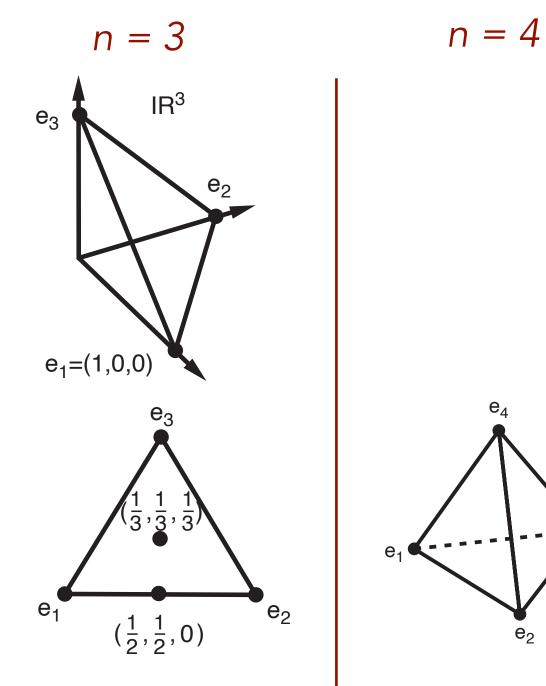
$$= x (1 - x) \left(f_1(x) - f_2(x) \right) = x (1 - x) \left((b - d) + (a - b - c + d) x \right)$$

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 and $f_2(x) = cx + d(1 - x)$

Replicator dynamics takes the form:

$$\dot{x} = x \left(f_1(x) - \bar{f}(x) \right) = x \left(f_1(x) - x f_1(x) - (1 - x) f_2(x) \right)$$

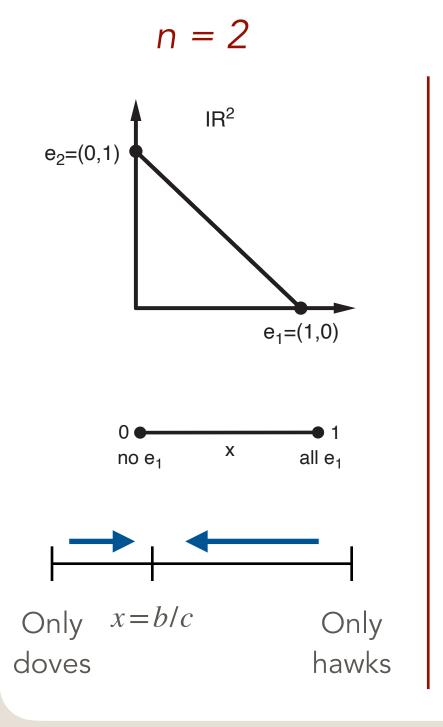
$$= x (1 - x) \left(f_1(x) - f_2(x) \right) = x (1 - x) \left((b - d) + (a - b - c + d) x \right)$$

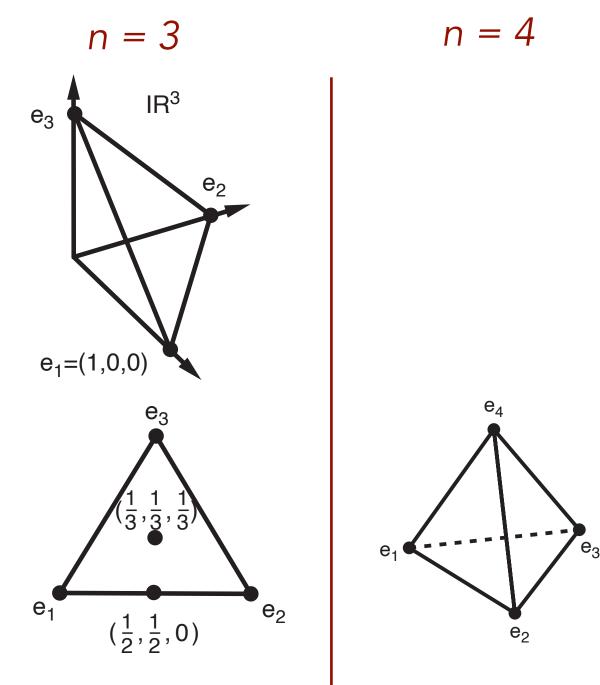
Fixed points: (1) Corners: x = 0, x = 1

Remark 1.6: On representing the unit simplex

Consider the replicator equation $\dot{x}_i = x_i (f_i(\mathbf{x}) - \bar{f}(\mathbf{x}))$. For a game with n strategies in total, this is, in principle, an *n*-dimensional system. However, we are only interested in those orbits on the unit simplex:

$$S_n = \left\{ \mathbf{z} \in \mathbb{R}^n : z_i \ge 0 \text{ for all } i \text{ and } \sum_{i=1}^n z_i = 1 \right\}$$





Remark 1.7: A classification of 2x2 games

To get some intuition, let us analyze the simplest non-trivial case: a symmetric game with two strategies:

	Action I	Action 2
Action 1	а	b
Action 2	С	d

We can represent the replicator equation as a 1-dim. system. Let x be the proportion of individuals who use Action 1, and 1-x is the proportion of individuals who use Action 2.

The fitnesses are

$$f_1(x) = ax + b(1 - x)$$
 and $f_2(x) = cx + d(1 - x)$

Replicator dynamics takes the form:

$$\dot{x} = x \left(f_1(x) - \bar{f}(x) \right) = x \left(f_1(x) - x f_1(x) - (1 - x) f_2(x) \right)$$

$$= x (1 - x) \left(f_1(x) - f_2(x) \right) = x (1 - x) \left((b - d) + (a - b - c + d) x \right)$$

Fixed points: (1) Corners: x = 0, x = 1

(2) Interior:
$$x = \frac{d-b}{a-b-c+d}$$
, if $x \in (0,1)$

Examples 1.8: Some 2x2 games

1. The hawk-dove game (with b=2, c=4)

	Hawk	Dove
Hawk	-1	2
Dove	0	1

Examples 1.8: Some 2x2 games

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			coexistence

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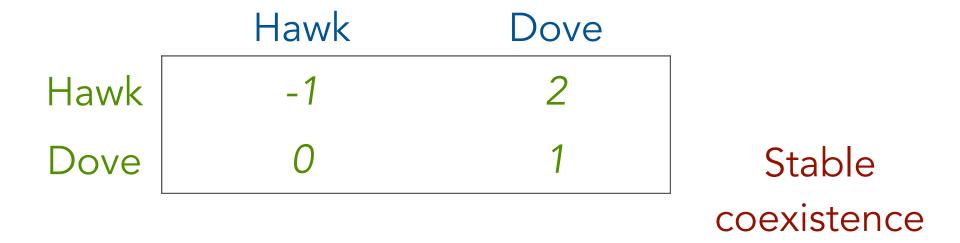
1. The hawk-dove game (with b=2, c=4)

$$\dot{x} = x(1-x)(1-2x)$$
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	Stag	Hare
Stag	10	0
Hare	7	7

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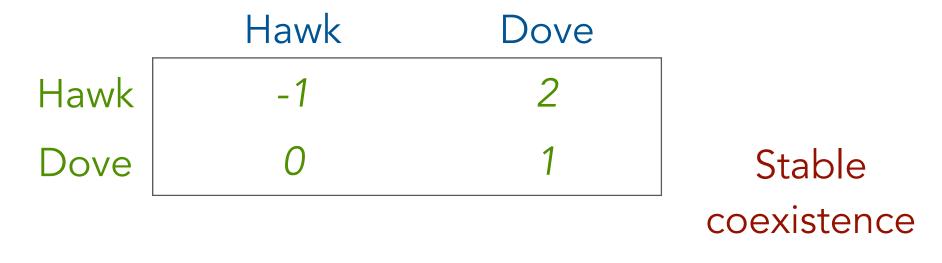
$$\dot{x} = x(1-x)(1-2x)$$
 Only doves $x = 1/2$ Only hawks $x = 0$

	Stag	Hare
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Hare	7	7

$$\dot{x} = x(1 - x)(10x - 7)$$

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 Only doves $x = 1/2$ Only hawks $x = 0$

	Stag	Hare
Stag	10	0
Hare	7	7

$$\dot{x} = x(1-x)(10x-7)$$
Only Hare
$$x = 7/10$$
Only Stag
$$x = 1$$

Examples 1.8: Some 2x2 games

1. The hawk-dove game (with b=2, c=4)

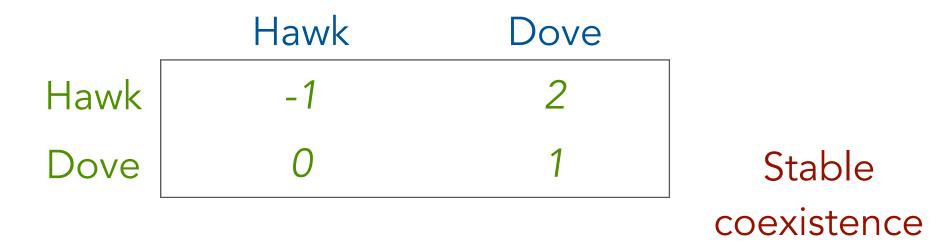


$$\dot{x} = x(1-x)(1-2x)$$
 Only doves $x = 1/2$ Only hawks $x = 0$

_	Stag	Hare	_
Stag	10	0	
Hare	7	7	
L			Bistability
$\dot{x} = x(1-x)(10x - x)$	· 7)	_	
	Only Hare	x =	7/10 Only Stag
	x = 0		x = 1

Examples 1.8: Some 2x2 games

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$$\dot{x} = x(1-x)(1-2x)$$
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2. Stag-hunt game (coordination game)

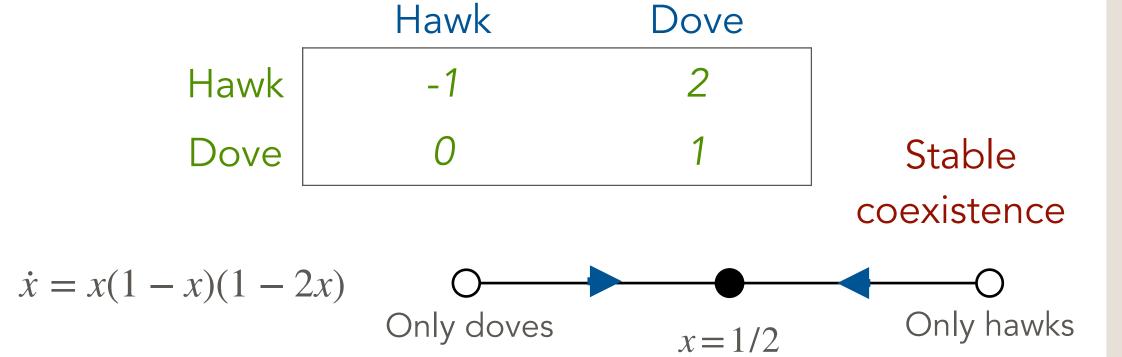
_	Stag	Hare	
Stag	10	0	
Hare	7	7	
L			Bistability
$\dot{x} = x(1-x)(10x - x)$	- 7)		
	Only Hare	x =	=7/10 Only Stag
	x = 0		x = 1

3. Prisoner's dilemma

	Cooperate	Defect
Cooperate	2	-1
Defect	3	0

Examples 1.8: Some 2x2 games

1. The hawk-dove game (with b=2, c=4)



2. Stag-hunt game (coordination game)

x = 0

	Stag	Hare	_
Stag	10	0	
Hare	7	7	
ı			Bistability
$\dot{x} = x(1-x)(10x - x)$	- 7)		
	Only Hare	x =	=7/10 Only Stag
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3. Prisoner's dilemma

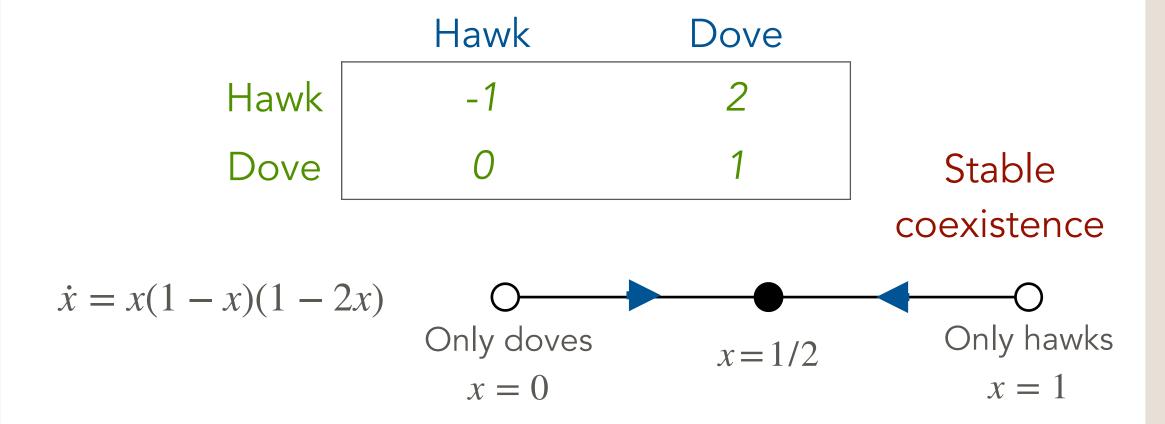
	Cooperate	Defect
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$$\dot{x} = x(1-x)(-1)$$

x = 1

Examples 1.8: Some 2x2 games

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$\dot{x} = x(1-x)(10x - $	7)	_	
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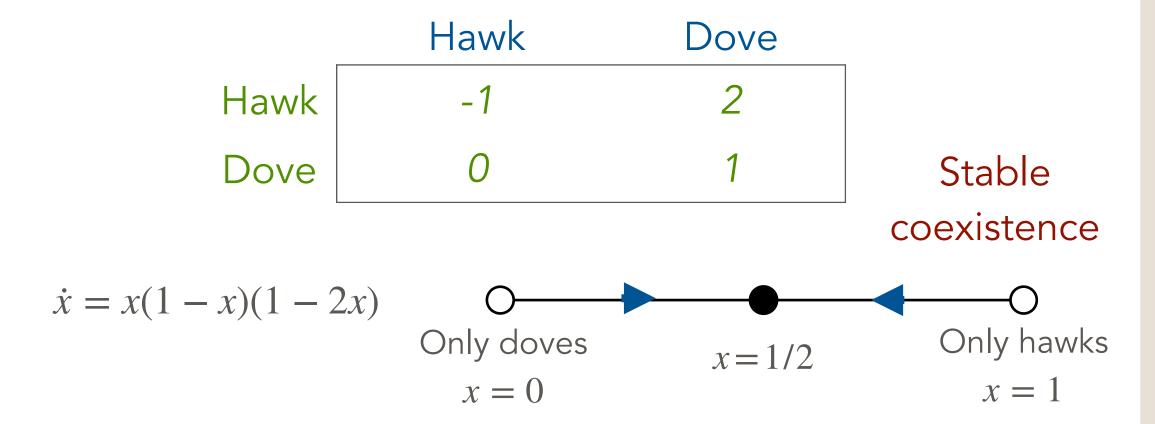
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$$\dot{x} = x(1-x)(-1)$$
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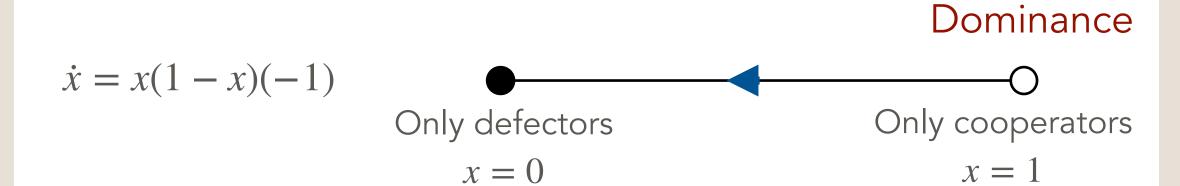


2. Stag-hunt game (coordination game)

_	Stag	Hare	
Stag	10	0	
Hare	7	7	
			Bistability
$\dot{x} = x(1-x)(10x - x)$	7)	_	
	Only Hare	x =	=7/10 Only Stag
	x = 0		x = 1

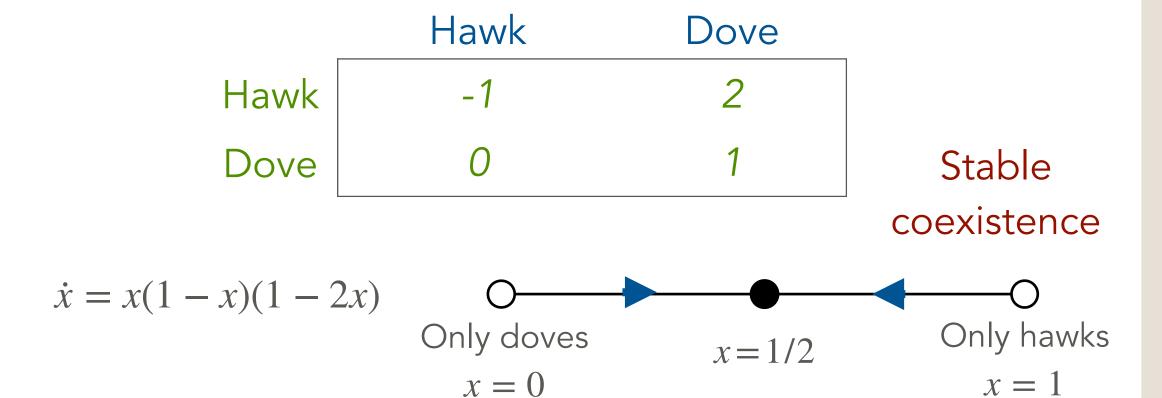
3. Prisoner's dilemma





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			Bistability
$\dot{x} = x(1-x)(10x - $	7)	_	
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	x = 0		x = 1

3. Prisoner's dilemma

	Cooperate	Defect
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Dominance

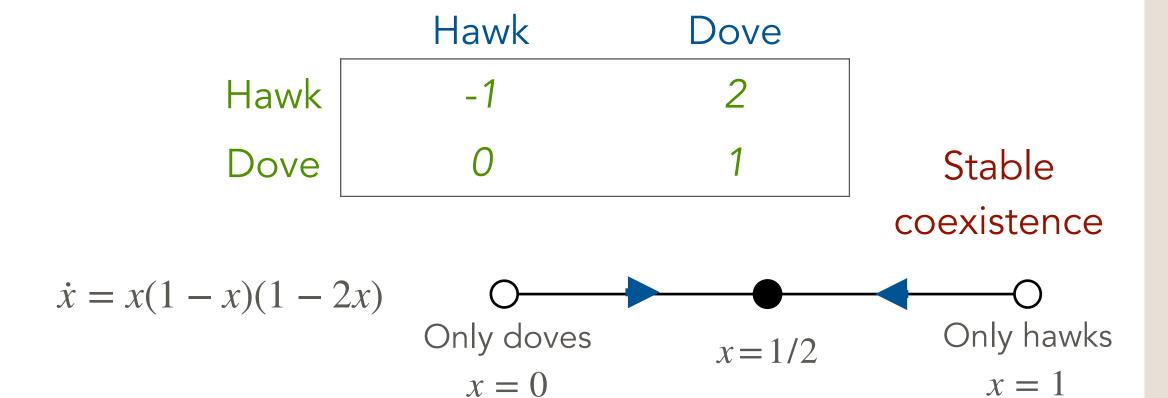
$$\dot{x} = x(1-x)(-1)$$
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Only cooperators
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4. A trivial game

	Action 1	Action 2
Action 1	3	1
Action 2	3	1

Examples 1.8: Some 2x2 games

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	x = 0		x = 1

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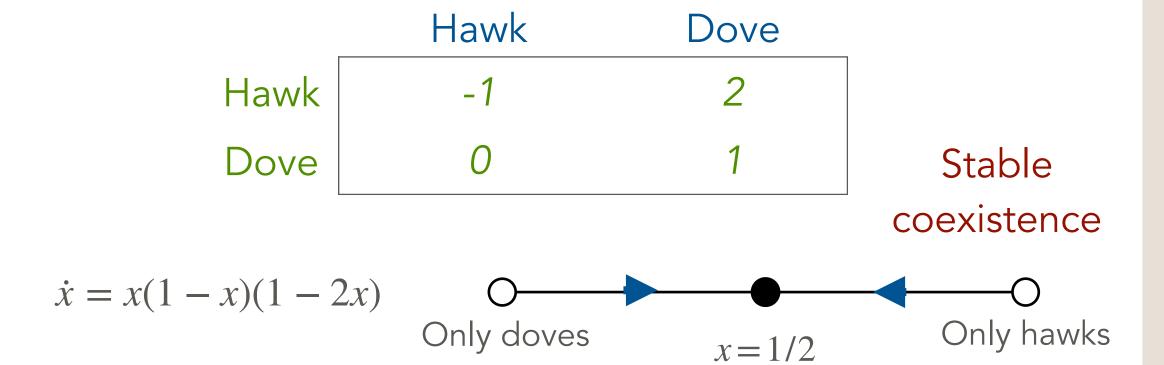
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	Action 1	Action 2
Action 1	3	1
Action 2	3	1

$$\dot{x} = x(1 - x) \cdot 0$$

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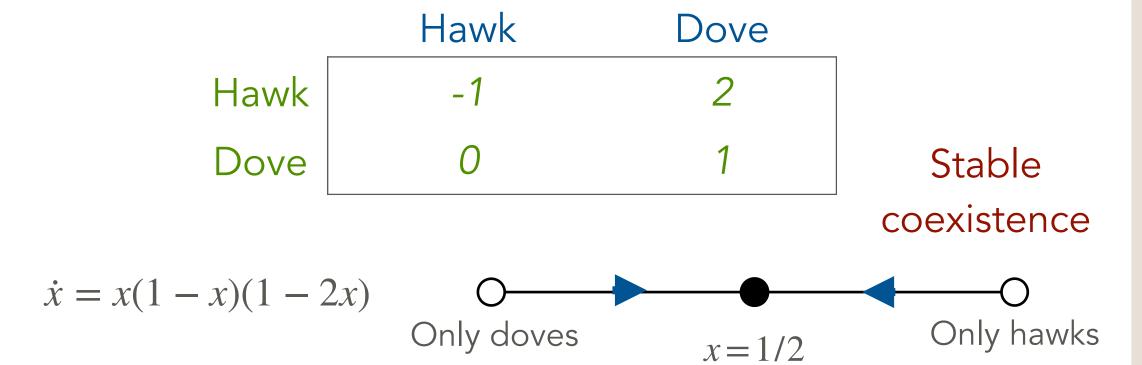
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Neutrality

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	x = 0			x = 1

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Neutrality

Qualitatively, these are all possible cases.

Example 1.9. A 3x3 game: The volunteer's timing dilemma

Consider the following variant of a so-called volunteer's dilemma. There are two players; at least one of them should volunteer to do a task that benefits both of them.

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1. Dynamics at the edges:

- No defectors ($x_D = 0$): Coexistence among cooperators and wait & see, $\mathbf{x}_{CW}^* = (1/4, 3/4, 0)$
- No wait&see ($x_W = 0$): Coexistence among cooperators and defectors, $\mathbf{x}_{CD}^* = (2/5, 0, 3/5)$
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2. Fixed point in the interior:

If $\dot{x}_i = x_i (f_i(\mathbf{x}) - \bar{f}(\mathbf{x}))$ has a fixed point with $x_i > 0 \ \forall i$, it must hold that $f_i(\mathbf{x}) = \bar{f}(\mathbf{x}) \ \forall i$.

Equivalently, it must hold that $f_1(\mathbf{x}) = f_2(\mathbf{x}) = f_3(\mathbf{x})$.

This is a simple linear system (and either has 0, 1, or infinitely many solutions).

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Equivalently, it must hold that $f_1(\mathbf{x}) = f_2(\mathbf{x}) = f_3(\mathbf{x})$.

This is a simple linear system (and either has 0, 1, or infinitely many solutions).

In our case, solution: $\mathbf{x}_{int}^* = (1/4, 3/16, 9/16)$.

Example 1.9. The volunteer's timing dilemma (continued)

3. Local stability analysis for the fixed points

Example 1.9. The volunteer's timing dilemma (continued)

- 3. Local stability analysis for the fixed points
 - For the equilibria on the edges, in each case it is true that the missing strategy can invade when rare.

Example 1.9. The volunteer's timing dilemma (continued)

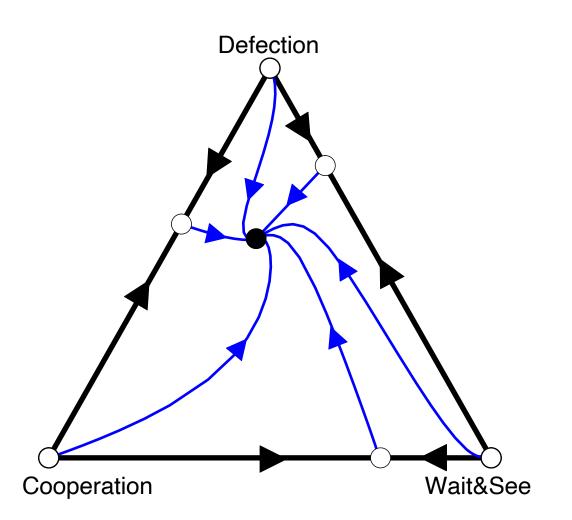
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 - For the equilibria on the edges, in each case it is true that the missing strategy can invade when rare.
 - The interior equilibrium is stable.

Example 1.9. The volunteer's timing dilemma (continued)

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4. Plotting some orbits numerically

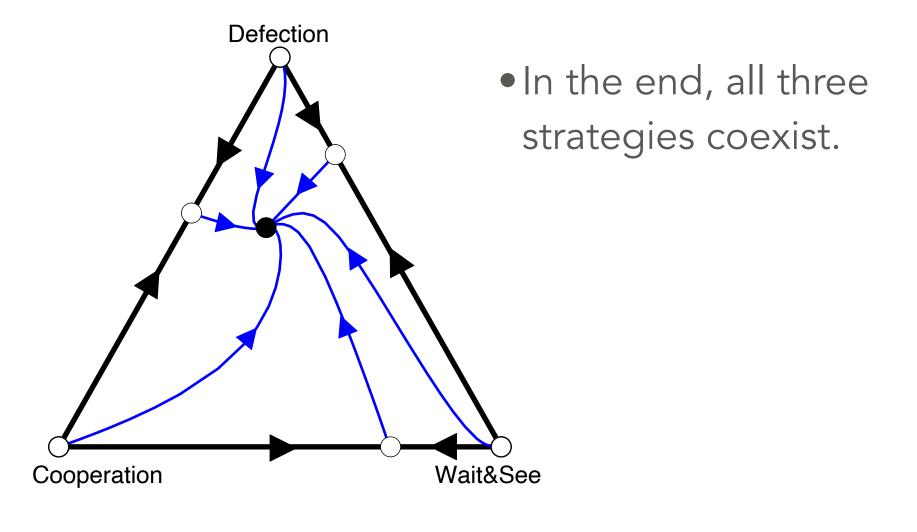


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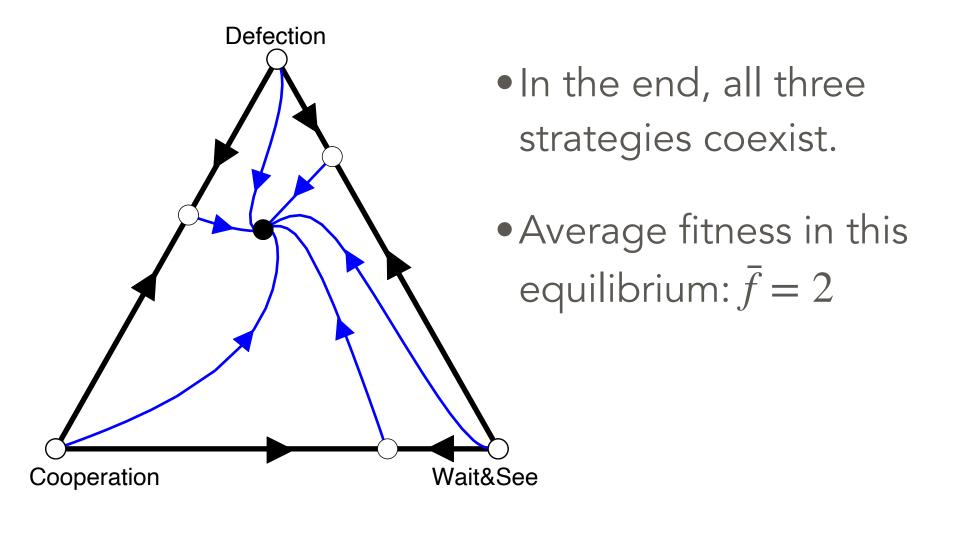


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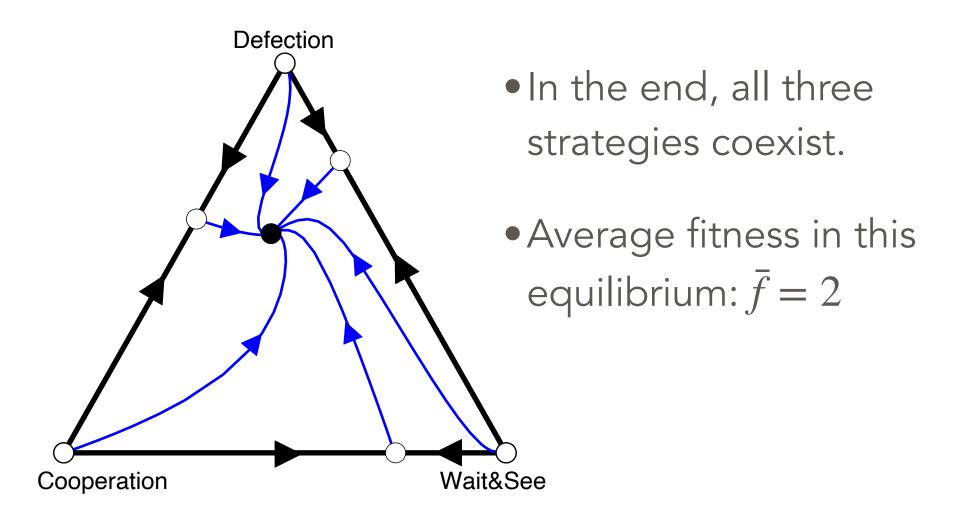


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Example 1.10. Rock Paper Scissors

Consider the following generalised version of rock paper scissors.

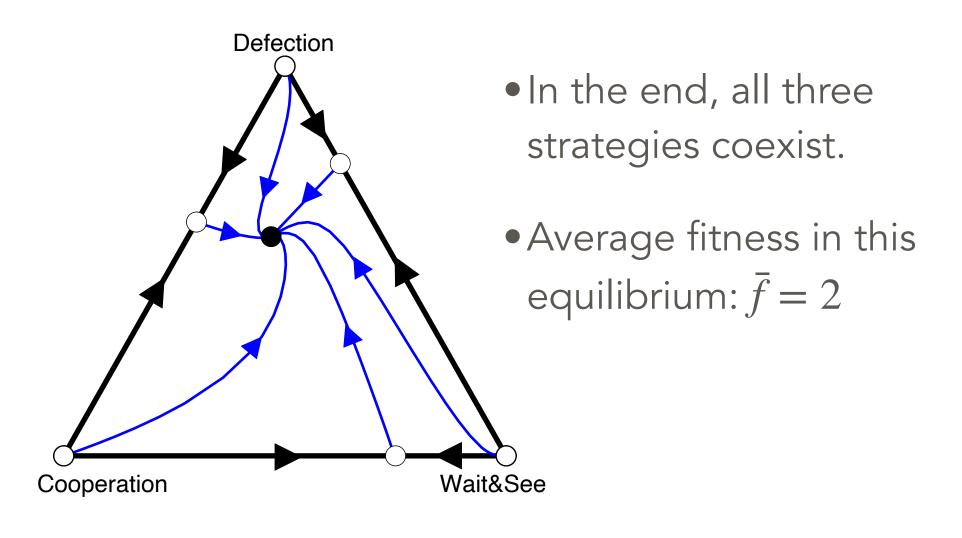
	Rock	Paper	Scissors
Rock	0	-a ₂	b ₃
Paper	b ₁	0	-a 3
Scissors	-a ₁	b_2	0

Example 1.9. The volunteer's timing dilemma (continued)

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	Rock	Paper	Scissors
Rock	0	-a ₂	<i>b</i> ₃
Paper	b ₁	0	-a 3
Scissors	-a ₁	b_2	0

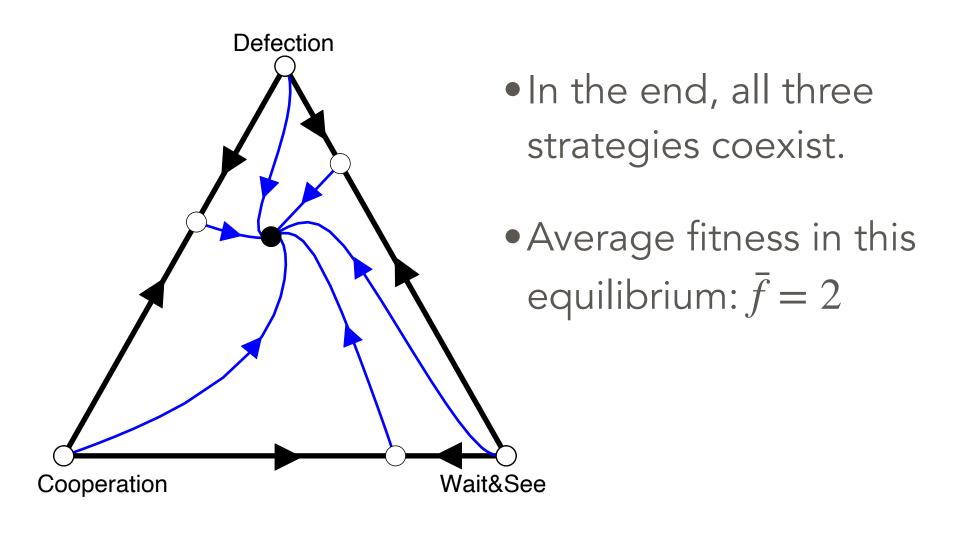
It turns out there are three possible dynamics.

Example 1.9. The volunteer's timing dilemma (continued)

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- For the equilibria on the edges, in each case it is true that the missing strategy can invade when rare.
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4. Plotting some orbits numerically



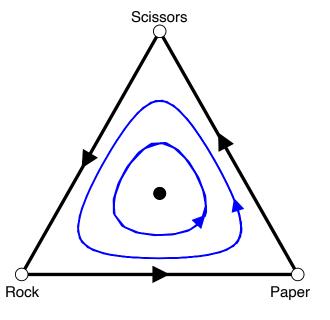
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	Rock	Paper	Scissors
Rock	0	-a ₂	b ₃
Paper	b ₁	0	-a ₃
Scissors	-a ₁	b_2	0

It turns out there are three possible dynamics.

$$a_1a_2a_3 = b_1b_2b_3$$

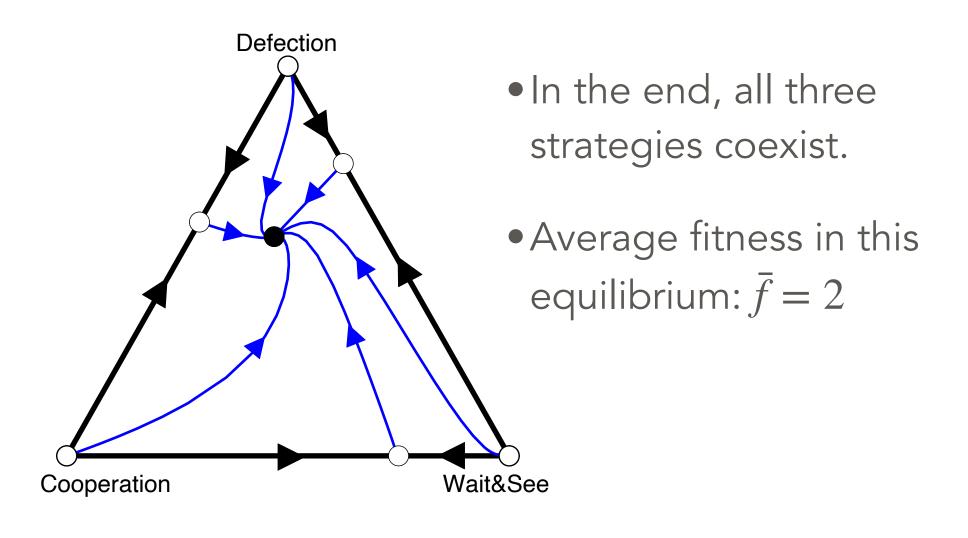


Example 1.9. The volunteer's timing dilemma (continued)

3. Local stability analysis for the fixed points

- For the equilibria on the edges, in each case it is true that the missing strategy can invade when rare.
- The interior equilibrium is stable.

4. Plotting some orbits numerically

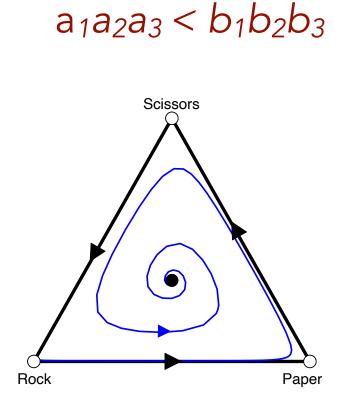


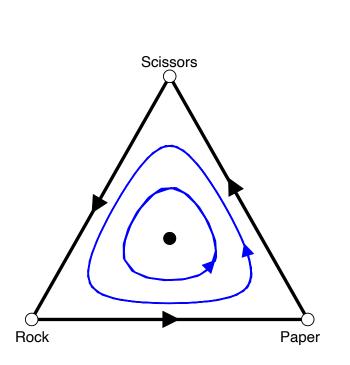
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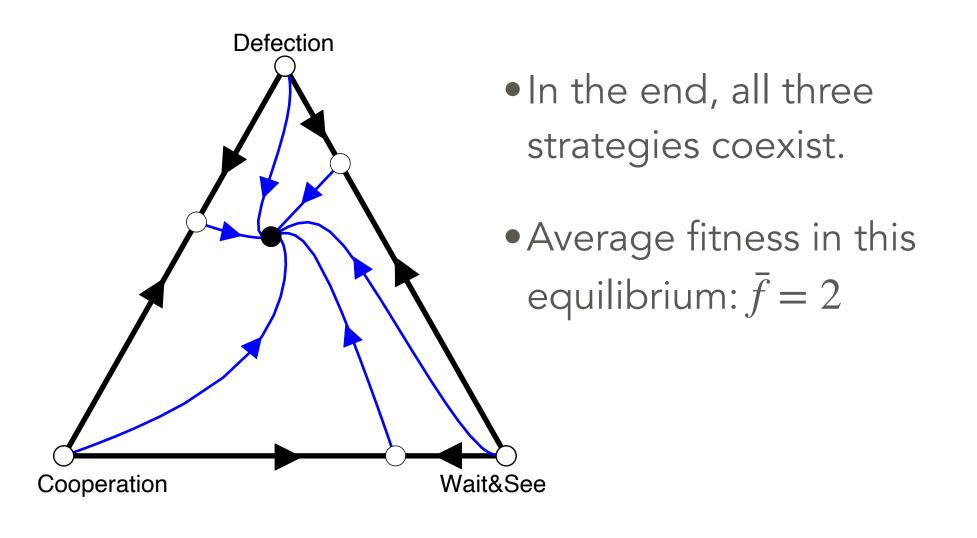
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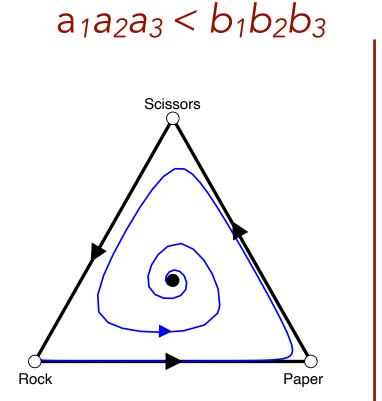


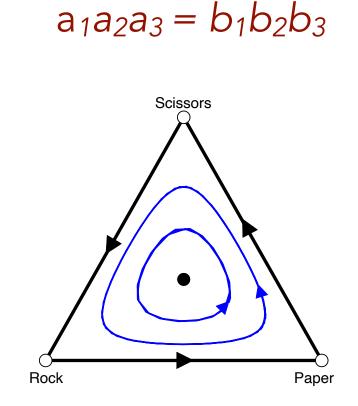
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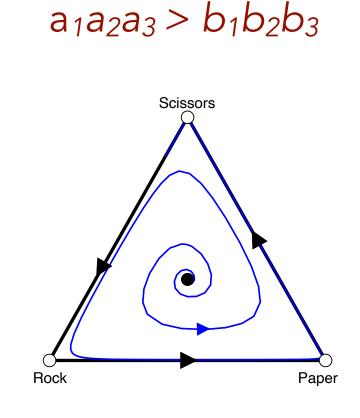
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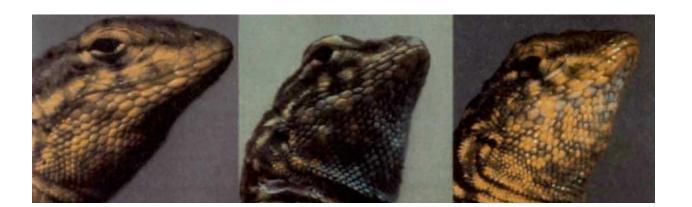
Evolutionary game theory: Non-transitive game in nature

Example 1.11. Non-transitive games in nature

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Three male morphs in side-blotched lizards:

- Males with orange throats defend large territories
- Males with blue throats defend smaller territories
- Males with yellow throats are sneakers without territory



The rock-paper-scissors game and the evolution of alternative male strategies

B. Sinervo & C. M. Lively

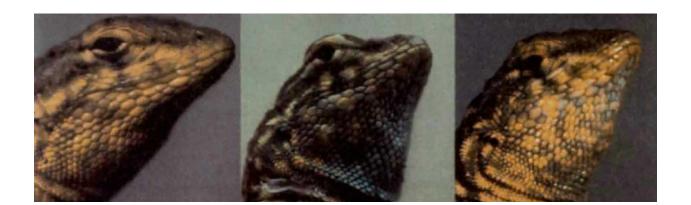
Department of Biology and Center for the Integrative Study of Animal Behavior, Indiana University, Bloomington, Indiana 47405, USA

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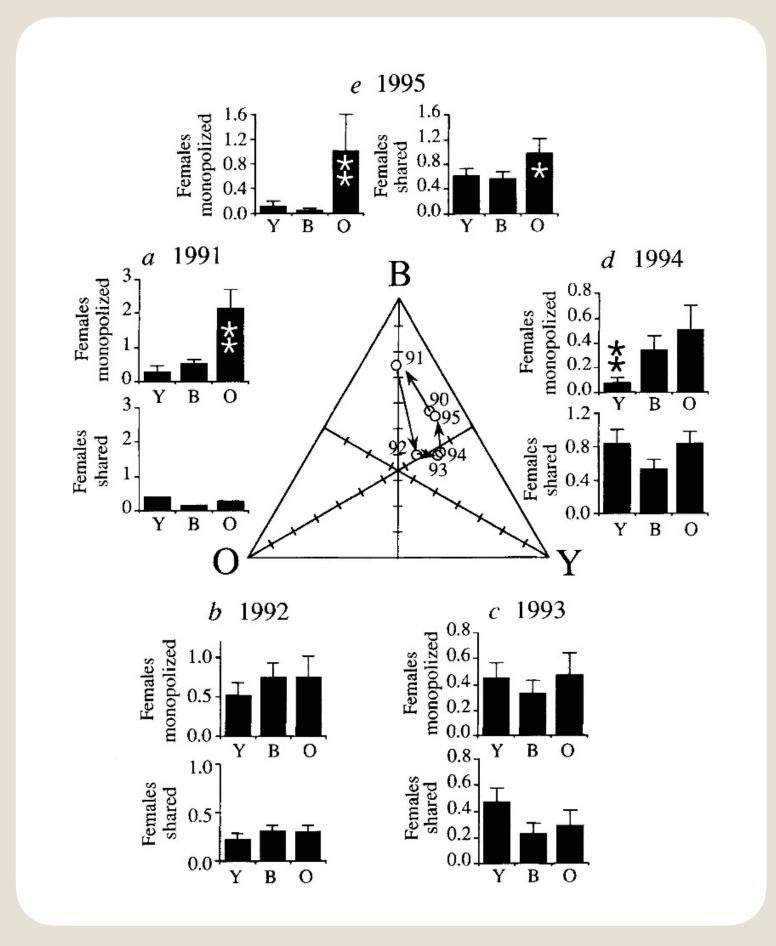
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Example 1.11. Non-transitive games in nature (continued)

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Three strains of E. Coli

- Colicin-producing strain (C)
- Sensitive strain (S)
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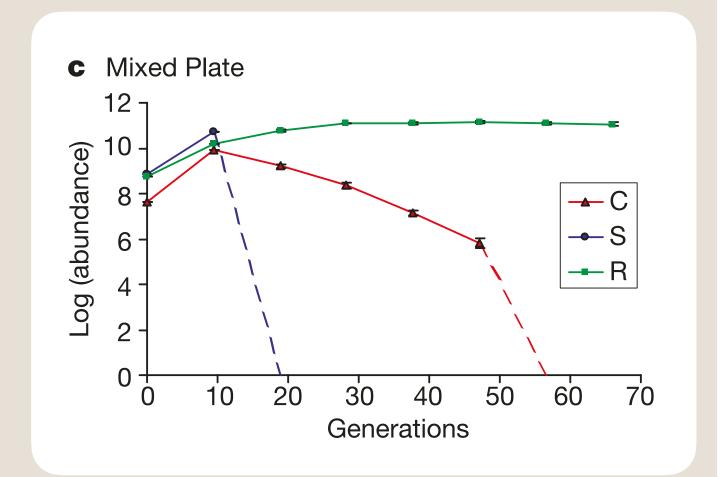
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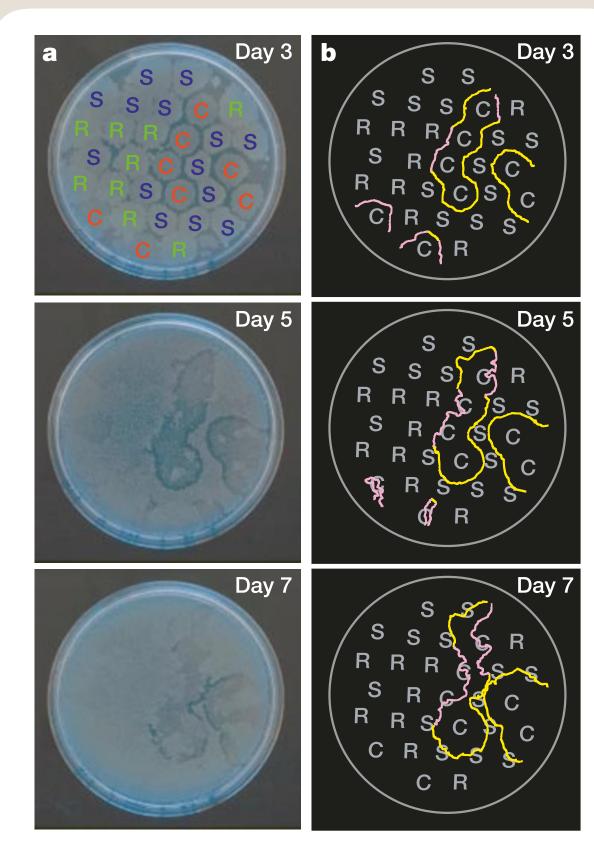
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In the examples we have seen so far, the outcome "predicted" by replicator dynamics often had a close relationship to the (symmetric) Nash equilibria of the game. This is not a coincidence; instead one can show the following results.

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Today, the replicator equation is one of the standard models of evolutionary game theory, for various reasons:

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- 4. There are beautiful connections to the concepts of classical game theory, without making any strong assumptions on the rationality of individuals. (How is this possible?)

Remark 2.14. Beyond replicator dynamics

Replicator dynamics might be both considered as a model of biological evolution, or of cultural evolution (imitation). However, it is also important to stress that replicator dynamics is one out of many evolutionary dynamics to consider. The optimal model depends on the applications one has in mind.

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To provide some intuition for how other models look like, I briefly discuss in the following the case of finite populations.

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• Everyone is equally likely to interact. As a result, if there are i individuals with strategy 1, the players' expected payoffs are given by:

$$\pi_1(i) = \frac{i-1}{N-1}a + \frac{N-i}{N-1}b$$

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• Setting $y_i := \varphi_i - \varphi_{i-1}$, we can write this as

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• Now, we use two different methods to sum up over all y_i . By its definition, we have

$$\sum_{i=1}^{N} y_i = (\varphi_1 - \varphi_0) + (\varphi_2 - \varphi_1) + \dots + (\varphi_N - \varphi_{N-1}) = \varphi_N = 1.$$

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- Let φ_i be the probability that eventually everyone adopts strategy 1, given that there are currently i players with this strategy.
- We can derive a recursion for ϕ_i

$$\varphi_i = T_i^+ \varphi_{i+1} + T_i^- \varphi_{i-1} + (1 - T_i^+ - T_i^-) \varphi_i$$
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with $\varphi_0 = 0$ and $\varphi_N = 1$.

• Setting $y_i := \varphi_i - \varphi_{i-1}$, we can write this as

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• In particular,

$$y_i = \prod_{k=1}^{i-1} \frac{T^-(k)}{T^+(k)} \varphi_1 \qquad (2.16.1)$$

• Now, we use two different methods to sum up over all y_i . By its definition, we have

$$\sum_{i=1}^{N} y_i = (\varphi_1 - \varphi_0) + (\varphi_2 - \varphi_1) + \dots + (\varphi_N - \varphi_{N-1}) = \varphi_N = 1.$$

On the other hand, using (2.16.1), we obtain:

$$\sum_{i=1}^{N} y_i = \varphi_1 \cdot \sum_{i=1}^{N} \prod_{k=1}^{i-1} \frac{T^-(k)}{T^+(k)} = \varphi_1 \cdot \left(1 + \sum_{i=1}^{N-1} \prod_{k=1}^{i} \frac{T^-(k)}{T^+(k)} \right)$$

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Because the two expressions need to coincide, we get

$$\varphi_1 = \frac{1}{1 + \sum_{i=1}^{N-1} \prod_{k=1}^{i} \frac{T^{-(k)}}{T^{+(k)}}}$$

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• In the special case that the game is a coordination game like stag-hunt (a > c, d > b), condition (2.17.1) is equivalent to $x^* < 1/3$, where

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This x^* is precisely the interior fixed point according to replicator dynamics.

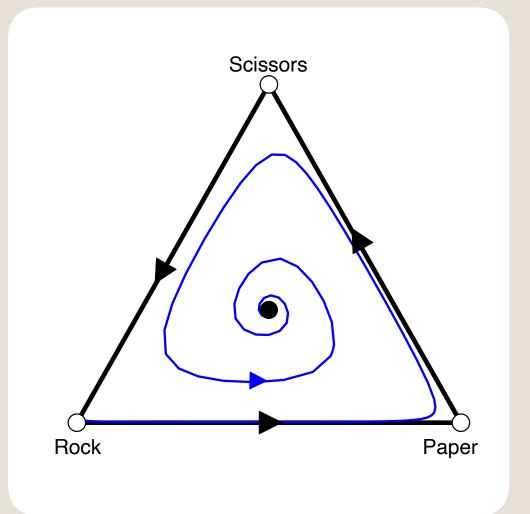


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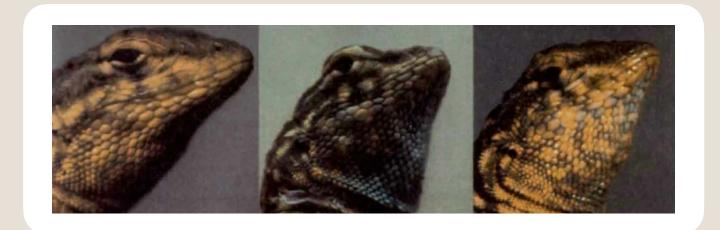


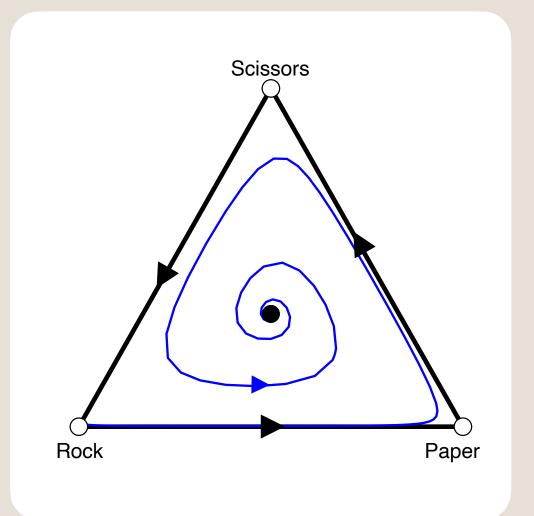


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- 2. Both dynamics have interesting mathematical properties, and they are well-connected to each other (and to the concepts of classical game theory; without making any a priori assumptions on the rationality of players).
- 3. Tomorrow, we will use such models of evolutionary dynamics to address one particular problem in evolutionary biology: why do individuals cooperate?

