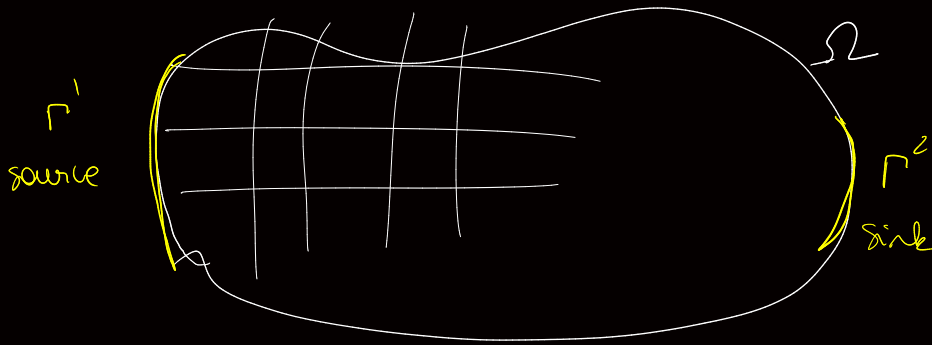


maximal flow and minimal cutset in FPP on $\mathbb{Z}^d$

I. Introduction of the model and definitions

$d \geq 2$ $(\mathbb{Z}^d, \mathbb{E}^d)$ G distribution on $(0, M)$ $M > 0$

$(c(e))_{e \in \mathbb{E}^d}$ i.i.d. family with law G



$c(e)$ = capacity of the edge e .
maximal amount of water that can flow through e / unit of time

definition: stream $f: \mathbb{E}^d \rightarrow \mathbb{R}^d$

$f(e)$ — $\|f(e)\|_2$: amount of water that flow through e

$\begin{matrix} \rightarrow \\ \leftarrow \end{matrix}$ $\frac{f(e)}{\|f(e)\|_2}$: $\begin{matrix} \rightarrow \\ \leftarrow \end{matrix}$
 e

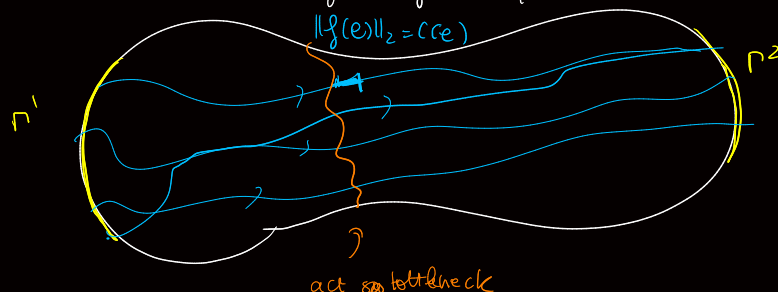
an admissible stream in $(\Omega, \Gamma^1, \Gamma^2)$:

- $f = 0$ outside Ω
- $\forall e \in \mathbb{E}^d \quad \|f(e)\|_2 \leq c(e)$
- the node law is respected everywhere outside Γ^1 and Γ^2

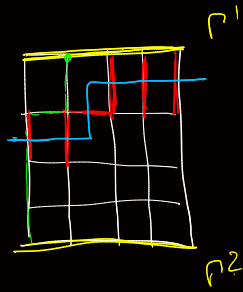


the flow of a stream $\text{flow}(f) = \sum_{e \in \Gamma^1 \nearrow \Gamma^2} \|f(e)\|_2$

maximal flow : $\phi(\Omega, \Gamma^1, \Gamma^2) = \sup \{ \text{flow}(f) : f \text{ admissible stream} \}$
if f is such that $\text{flow}(f) = \phi(\Omega, \Gamma^1, \Gamma^2)$ we call f a maximal stream.



Def. [cutset in Ω from Γ^1 to Γ^2] $E \subset \mathbb{E}^d$ s.t. $\forall$ path γ from Γ^1 to Γ^2 in Ω , we have $E \cap \gamma \neq \emptyset$



capacity of a cutset E : $\text{cap}(E) = \sum_{e \in E} c(e)$.

max-flow min-cut theorem: $\phi(\Omega, \Gamma^1, \Gamma^2) = \inf \{ \text{cap}(E) : E \text{ cutset} \}$

admissible stream

$\downarrow$

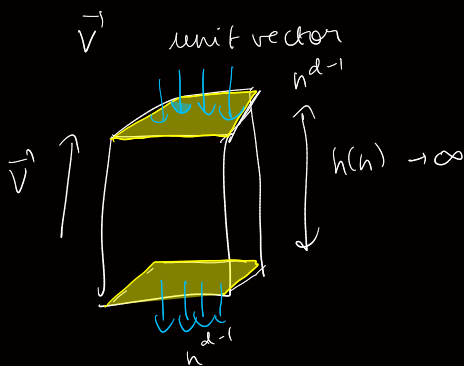
flow(f)

cutset

$\downarrow$

$\text{cap}(E)$

II. Maximal flow in a given direction



$$\phi_n = \phi(\square_{n^{d-1}}, \square, \square)$$

$$\frac{\phi_n}{n^{d-1}}$$

Thm (Kesten 87)

$$\frac{\phi_n}{n^{d-1}} \rightarrow$$

$$\mathcal{V}_G(\vec{v})$$

$\wedge$
flow constant

$d=3$, $\vec{v} \parallel$ to axis

$h(n)$ polynomial in n

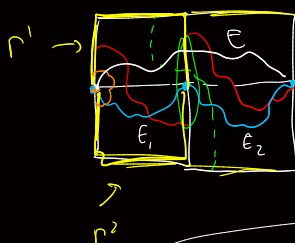
Thm (Borovik & Thérèse 10)

$$\frac{\phi_n}{n^{d-1}} \rightarrow$$

$$\mathcal{V}_G(\vec{v})$$

any d .
any $\vec{v}$

$\frac{h(n)}{n} \rightarrow 0$
 $h(n) \rightarrow \infty$



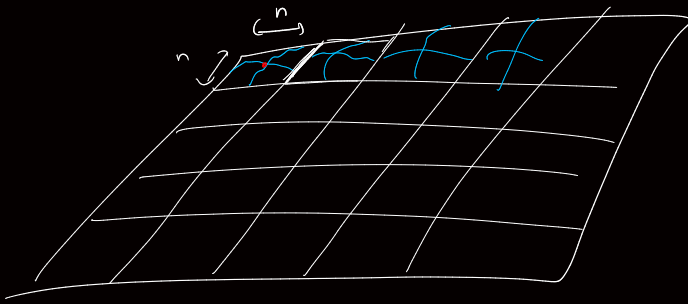
$$\tau_n = \phi(\square_{n^{d-1}}, \square, \square)$$

$$\text{cap}(E) \leq \text{cap}(E_1 \cup E_2) = \text{cap}(E_1) + \text{cap}(E_2)$$

$$\frac{\tau_n}{n^{d-1}}$$

$$N \gg n$$

$$|\mathcal{I}| \approx \left(\frac{N}{n}\right)^{d-1}$$



$$\tau_N \leq \sum_{i \in \mathcal{I}} \tau_n(i)$$

$$\mathbb{E}(\tau_N) \leq \left(\frac{N}{n}\right)^{d-1} \mathbb{E}(\tau_n)$$

$$\frac{1}{N^{d-1}} \mathbb{E}(\tau_N) \leq \frac{1}{n^{d-1}} \mathbb{E}(\tau_n)$$

$$\limsup_{N \rightarrow \infty} \frac{1}{N^{d-1}} \mathbb{E}(\tau_N) \leq \lim_{n \rightarrow \infty} \frac{1}{n^{d-1}} \mathbb{E}(\tau_n)$$

$$\Rightarrow \lim_{n \rightarrow \infty} \frac{1}{n^{d-1}} \mathbb{E}(\tau_n) = \lim_{n \rightarrow \infty} \frac{\mathbb{I}_n}{n^{d-1}} = V_G(\vec{v}) \quad \text{a.s.}$$

Concentration argument.

$V_G(\vec{v})$: maximal amount of water that can flow in direction $\vec{v}$ asymptotically



$$h(n) \leq n$$

$$h(n)n^{d-2}$$

$$E \cup \frac{1}{n} \frac{1}{n}$$

$$\phi_n \leq \tau_n$$

$$\tau_n \leq \phi_n + \underbrace{\left\{ \right\}}_{h(n)n^{d-2}}$$

$$\lim_{n \rightarrow \infty} \frac{\tau_n}{n^{d-1}} = \lim_{n \rightarrow \infty} \frac{\phi_n}{n^{d-1}}$$

μ distributed on $[0, n]$

$(X_i)_{i \leq n}$ i.i.d distributed given μ

$$\text{LLN} \quad \frac{1}{n} \sum_{i=1}^n X_i \rightarrow \mathbb{E}(X_1)$$

$$P\left[\frac{1}{n} \sum_{i=1}^n X_i \geq \mathbb{E}(X_1) + \varepsilon\right] \approx e^{-\frac{n}{\varepsilon} I(\varepsilon)}$$

speed of large deviation

cost

Large deviations for $\bar{c}_n$ 

Lower large deviations $P(\frac{\bar{c}_n}{n^{d-1}} \leq \nu_G(\bar{v}) - \varepsilon) \approx e^{-n^{d-1} J_G(\varepsilon)}$

Rossignol & Thérèse
2010



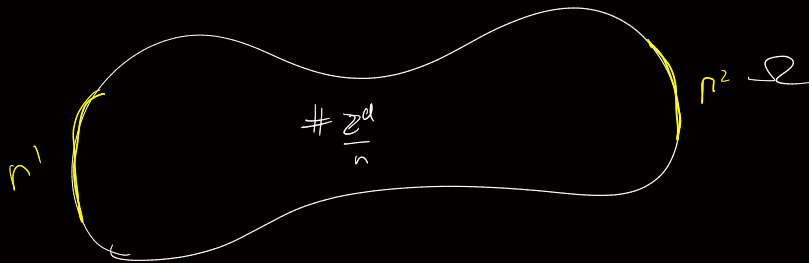
$$\varepsilon (\nu_G(\bar{v}) - \varepsilon) n^{d-1}$$

Upper large deviations: $P(\frac{\bar{c}_n}{n^{d-1}} \geq \nu_G(\bar{v}) + \varepsilon)$ of volumic order



$$\limsup_{n \rightarrow \infty} \frac{1}{n^d} \log P(\quad) > 0 \quad \approx e^{-n^d}$$

III. maximal flow in a general domain



$$\phi_n(\Omega, r^1, r^2)$$

Thm: (Corollary of Theorem 11) $\phi_\Omega \geq 0$

$$\phi_n(\Omega, r^1, r^2) \rightarrow \phi_\Omega \text{ a.s.}$$

$E_n^{\min}$ minimal cut set

$E_n^{\min} \xrightarrow{n^{d-1}} \text{continuous minimal cut set}$

$f_n^{\max}$ maximal stream

$f_n^{\max} \xrightarrow{\quad} \text{continuous maximal stream}$

We want to encode f_n into a measure

$$\bar{\mu}_n(f_n) = \sum_{e \in \mathbb{E}_n^d} \underbrace{f_n(e)}_{\in \mathbb{R}^d} \underbrace{\delta_{C(e)}}_{\text{dirac mass at the center of } e} \in \mathcal{M}(\mathbb{R}^d)^d$$



vector field

$\vec{\sigma} = 0$ outside Ω .

$\text{div } \vec{\sigma} = 0$ in $\Omega \Leftrightarrow$ node law

ULD

$$\lambda > \phi_{-2} \quad \mathbb{P}(\phi_n \geq \lambda n^{d-1})$$

$$\text{Thm (D, Theoret 20)} \quad \forall A \xrightarrow{\text{Borelian set}} \mathcal{B}(\mathcal{M}(\mathbb{R}^d)^d) \quad , \quad I \geq 0$$

$$-\inf_{\vec{A}} I \leq \liminf_{n \rightarrow \infty} \frac{1}{n^d} \log \mathbb{P}(\cdot) \leq \limsup_{n \rightarrow \infty} \frac{1}{n^d} \log \mathbb{P}(\exists f_n \text{ s.t. } \vec{\mu}_n(f_n) \in A) \leq \inf_A I$$

$$\mathbb{P}(\exists f_n \text{ s.t. } \vec{\mu}_n(f_n) \approx \vec{\sigma} \mathbb{Z}^d) \approx e^{-\underbrace{I(\vec{\sigma})}_{I(\vec{\sigma}) = \int_{\mathbb{R}^d} \frac{1}{I(\vec{\sigma}(x))} d\varphi^d(x)} n^d}$$

$$J(\lambda) = \inf \{ I(\vec{\sigma}) : \underline{\text{flow}}(\vec{\sigma}) \geq \lambda \}$$

$$\frac{\text{flow}(f_n)}{n^{d-1}} \geq 1$$

$$\vec{\mu}_n(f_n) \approx \vec{\sigma} \mathbb{Z}^d$$

$$\mathbb{P}(\phi_n \geq \lambda n^{d-1}) = \mathbb{P}(\exists f_n \text{ admissible s.t. } \text{flow}(f_n) \geq \lambda n^{d-1})$$

$$\approx \mathbb{P}(\exists f_n \text{ s.t. } \vec{\mu}_n(f_n) \in \{ \vec{\sigma} \mathbb{Z}^d : \underline{\text{flow}}(\vec{\sigma}) \geq 1 \})$$

$$\approx e^{-n^d J(\lambda)}$$

$$\mathbb{P}(\phi_n \leq \lambda n^{d-1}) \quad \lambda < \phi_{-2}$$