

MLSI: THEORY, EXAMPLES AND CONSEQUENCES LECTURE 3

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1 DISCRETE CURVATURE AND CONJECTURES

2 SLC AND MATROID BASES EXCHANGE

BOCHNER-BAKRY-ÉMERY CRITERION

The following 2nd derivative criterion is also useful, and inspires a notion of *discrete curvature*, mimicking the Ricci curvature.

PROPOSITION (BOCHNER'46, LICHNÉROWICZ'58)

$$\inf_f \frac{\mathcal{E}(f, f)}{\text{Var}_\pi f} =: \lambda = \inf_f \frac{\mathcal{E}(-Lf, f)}{\mathcal{E}(f, f)} =: \mu_P =: \mu.$$

Proof. $\lambda \leq \mu$: Use Cauchy-Schwartz. Indeed, w/ $\mathbb{E}_\pi f = 0$,

$$\begin{aligned} \mathcal{E}(f, f) = \mathbb{E}(f(-Lf)) &\leq \left(\text{Var}_\pi f\right)^{1/2} \left(\mathbb{E}(-Lf)^2\right)^{1/2} \\ &\leq \frac{1}{\sqrt{\lambda}} (\mathcal{E}(f, f))^{1/2} \left(\mathbb{E}(-Lf)^2\right)^{1/2} \\ &= \frac{1}{\sqrt{\lambda}} (\mathcal{E}(f, f))^{1/2} \left(\mathcal{E}(-Lf, f)\right)^{1/2}. \end{aligned}$$

BOCHNER-BAKRY-ÉMERY CRITERION (CONTD.)

Other direction, $\lambda \geq \mu$: 2nd derivative and integration. Starting with

$$\frac{d}{dt} \text{Var}(H_t f) = -2\mathcal{E}(H_t f, H_t f) \text{ and } \frac{d}{dt} \mathcal{E}(H_t f, H_t f) = -2\mathcal{E}(H_t f, -L H_t f).$$

Observe that as $t \rightarrow \infty$, $H_t f \rightarrow \mathbb{E}_\pi f$, and so $\mathcal{E}(H_t f, H_t f) \rightarrow 0$. Hence,

$$\begin{aligned} \mathcal{E}(f, f) &= - \int_0^\infty \frac{d}{dt} \mathcal{E}(H_t f, H_t f) dt = 2 \int_0^\infty \mathcal{E}(H_t f, -L(H_t f)) dt \\ &= 2 \int_0^\infty \mathcal{E}(-L^*(H_t f), H_t f) dt \geq 2\mu_{P^*} \int_0^\infty \mathcal{E}(H_t f, H_t f) dt \\ &= -\mu_{P^*} \int_0^\infty \frac{d}{dt} \text{Var}(H_t f) dt = \mu_{P^*} \text{Var} f, \end{aligned}$$

implying that $\lambda_P = \lambda_{P^*} \geq \mu_P$.

MLSI AND BAKRY-ÉMERY CRITERION

What about repeating the above for the entropy constant?

PROPOSITION (BAKRY-ÉMERY'85)

$$\inf_{f>0} \frac{\mathcal{E}(f, \log f)}{\text{Ent}_\pi f} =: \alpha \geq \inf_{f>0} \frac{u(f)}{\mathcal{E}(f, \log f)},$$

where $u(f) = \mathcal{E}(-Lf, \log f) + \mathcal{E}(f, (-Lf)/f)$.

Some relevant references include:

[Boudou, Caputo, Dai Pra, Posta, 2005]. “Spectral gap estimates for interacting particle systems via a Bochner type identity,” J. Funct. Analysis.

[M. Erbar, J. Maas, P.T. 2015]. “Discrete Curvature Bounds for Bernoulli-Laplace and Random Transposition Models,” Annales Fac. Sci. Toulouse.

Bochner formula (iterated gradient Γ_2 criterion) :
application: discrete Buser ...

BAKRY-ÉMERY, ... , SCHMUCKENSCHLÄGER,...

$G = (V, E)$, and Graph Laplacian $\Delta = -(D - A)$

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DEFINITION [BOCHNER w/ PARAMETER K]

Curvature of G is at least K , if $\forall f : V \rightarrow \mathbb{R}$, and $\forall x \in V$,

$$\Delta \Gamma(f, f)(x) - 2\Gamma(f, \Delta f)(x) \geq K \Gamma(f, f)(x),$$

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where

$$\Delta f(x) = \sum_{y:(x,y) \in E} (f(y) - f(x)).$$

And given $f, g : V \rightarrow \mathbb{R}$, also define:

$$\Gamma(f, g)(x) = \frac{1}{2} \sum_{y:(x,y) \in E} (f(x) - f(y))(g(x) - g(y)).$$

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Note: $\Gamma(f)(x) := \Gamma(f, f)(x) = \frac{1}{2} \sum_{y:(x,y) \in E} (f(x) - f(y))^2 =: |\nabla f(x)|^2$.

Γ_2 CALCULUS

Basically, using the iterated gradient:

$$2\Gamma_2(f, g) = \Delta\Gamma(f, g) - \Gamma(f, \Delta g) - \Gamma(\Delta f, g).$$

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EQUIVALENTLY ...

Curvature of G is at least K , if $\forall f : V \rightarrow \mathbb{R}$, and $\forall x \in V$,

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 2\Gamma_2(f)(x) &= \Delta\Gamma(f)(x) - 2\Gamma(f, \Delta f)(x) \\
 &= \sum_{v \sim x} \Gamma(f)(v) - d(x)\Gamma(f)(x) - \sum_{v \sim x} f(v)(\Delta f(v) - \Delta f(x)) \\
 &= \left(\sum_{v \sim x} f(v) \right)^2 - \frac{d(x)}{2} \sum_{v \sim x} f^2(v) + \sum_{u \sim v \sim x} \frac{f^2(u) - 4f(u)f(v) + 3f^2(v)}{2} \\
 &= \left(\sum_{v \sim x} f(v) \right)^2 - \sum_{v \sim x} \frac{d(x) + d(v)}{2} f^2(v) + \frac{1}{2} \sum_{u \sim v \sim x} ((f(u) - 2f(v))^2.
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 \end{aligned} \tag{2.1}$$

For d -regular, further simplifies to:

$$2\Gamma_2(f)(x) = d \sum_{v \sim x} f^2(v) + \left(\sum_{v \sim x} f(v)\right)^2 + \sum_{v \sim x} \sum_{u \sim v} \left(\frac{f^2(u)}{2} - 2f(u)f(v)\right). \tag{2.2}$$

CHEEGER IS TIGHT IF CURVATURE ≥ 0

$h := h(G) := \min_{A \subset V} |\partial A|/|A|$, the Cheeger const, of G ; $\lambda > 0$: $\text{gap}(G)$.

THEOREM (A-M, L-S, J-S 80's)

$$\frac{h^2}{c_1 \max \deg(G)} \leq \lambda \leq c_2 h.$$

THEOREM (KLARTAG-KOZMA-RALLI-T. '14)

1. Bochner w/ parameter K implies: for any subset $A \subset V$,

$$|\partial A| \geq \frac{1}{2} \min \left\{ \sqrt{\lambda}, \frac{\lambda}{\sqrt{2|K|}} \right\} |A| \left(1 - \frac{|A|}{|V|} \right).$$

2. Bochner w/ parameter $K \geq 0$ implies $\lambda \geq K$, and hence:

$$\lambda \leq 16h^2.$$

3. Abelian Cayley graphs satisfy Bochner w/ parameter $K \geq 0$.

CD(K, ∞) WITH $K \geq 0$: EXAMPLES

- 1. Complete graph K_n : Curvature = $1 + n/2$.
- 2. Discrete n -cube Q_n : Curvature = 2.
- 3. Symmetric group $S(n)$ with the transposition metric: Curvature = 2
- 4. **General Proposition.** If T is the maximum number of triangles containing any edge, then $K \leq 2 + T/2$.
- 5. **Abelian necessary!** On S_n , the (left) Cayley graph generated by $(12), (12\dots n) \pm 1$, the Cheeger constant is $c_1 n^{-2}$, while the spectral gap is $c_2 n^{-3}$, with $c_1, c_2 > 0$, independent of n .

CHEEGER, $CD(K, \infty)$ WITH $K \geq 0$

- Qn. 1 Can we characterize the graphs or Markov kernels which satisfy non-negative curvature?
- Qn. 2 Can we characterize the graphs for which the Cheeger inequality is tight?
- Qn. 3 More examples? Non-crossing partition lattice $NC(n)$?

OLLIVIER-VILLANI

“ In positive curvature, balls are closer than their centers.”

Coarse Ricci: Take two small balls and compute the transportation distance between them. If it is smaller than the distance between the centers of the balls, then coarse Ricci is *positive*.

$$W_1(\mu_x, \mu_y) =: (1 - \kappa(x, y))d(x, y),$$

Examples.

1. n -cube: μ_x uniform on the $n+1$ neighbors of x (including itself). For x, y , neighbors, $\kappa(x, y) = 2/(n+1)$.
2. S_n with transpositions: For σ, τ differing in a transposition, $\kappa(x, y) = 1/\binom{n}{2}$.

(COARSE) RICCI OF HYPERCUBE

PROPOSITION (FOLKLORE)

The n -cube has coarse Ricci $= 2/(n + 1)$.

Proof sketch.

- (i) Consider the *lazy* random walk on the n -cube.
- (ii) Use a simple (path) coupling argument to show that two copies of the chain started at *neighboring vertices* $x, y \in \{0, 1\}^n$ can be “coupled” with probability at least $2/(n + 1)$ in one step.

(COARSE) RICCI OF TRANSPOSITION GRAPH

PROPOSITION

S_n with transpositions has coarse Ricci = $1/\binom{n}{2}$.

Proof sketch.

Lower bound:

- (i) Consider the *lazy* random transposition chain on S_n .
- (ii) Use a simple (path) coupling argument to show that two copies of the chain started at $\sigma, \tau \in S_n$ with $d(\sigma, \tau) = 1$, can be coupled with probability at least $1/\binom{n}{2}$ in one step.

Upper bound: Follows by showing that there is no better coupling – use (Kantorovich's) dual formulation of W_1 :

$$W_1(\nu, \mu) = \sup_{f: 1\text{-Lip}} \left(E_\nu f - E_\mu f \right).$$

PERES-TETALI “CONJECTURE”!

Does *coarse Ricci* at least $\kappa > 0$ imply the MLSI constant $\alpha \geq \kappa$?

Something weaker is known: namely that a W_1 transport-Entropy inequality follows, thanks to (Marton??, Eldan, Lehec, Lee'16)

MLSI FOR BASES EXCHANGE WALK

I am thankful to Heng Guo for letting me borrow some of his slides in this section.

MATROID

A matroid $\mathcal{M} = (E, \mathcal{I})$ consists of a finite ground set E and a collection \mathcal{I} of subsets of E (independent sets) such that:

- $\emptyset \in \mathcal{I}$;
- if $S \in \mathcal{I}$, $T \subseteq S$, then $T \in \mathcal{I}$ (**downward closed**);
- if $S, T \in \mathcal{I}$ and $|S| > |T|$, then there exists an element $i \in S \setminus T$ such that $T \cup \{i\} \in \mathcal{I}$ (**augment axiom**).

*Maximum independent sets are the **bases**.*

For any two bases, there is a sequence of exchanges of ground set elements that take one basis to the other.

Let $n = |E|$ and r be the **rank**, namely the size of any basis.

BASES-EXCHANGE WALK

The following Markov chain $P_{\text{BX},\pi}$ converges to a “homogeneous SLC” π :

- 1 **remove** an element uniformly at random from the current basis (call the resulting set S);
- 2 **add** $i \notin S$ with probability proportional to $\pi(S \cup \{i\})$.

The implementation of the second step may be non-trivial.

MLSI FOR MATROID BASIS EXCHANGE

THEOREM (MARY CRYAN-HENG GUO-GIORGOS MOUSA)

For any $f : \Omega \rightarrow \mathbb{R}_{\geq 0}$,

$$\mathcal{E}_{P_{\text{BX},\pi}}(f, \log f) \geq \frac{1}{r} \cdot \text{Ent}_{\pi}(f),$$

where r is the rank of the matroid.

MANY OPEN PROBLEMS REMAIN!

- Sampling spanning trees of G with **no broken circuit** - the same as evaluation $T_G(1, 0)$ of the Tutte polynomial, leading coefficient $|a_1[\chi_G(\lambda)]|$ of the chromatic polynomial (up to sign), number of maximum G -parking functions, sampling acyclic orientations *with a unique sink* ...
- Sampling all acyclic orientations of G
- Negative correlation for random forests of G
- (More generally) characterization of matroids with negative correlation for random bases.

STRONGLY LOG-CONCAVE POLYNOMIALS

LOG-CONCAVE POLYNOMIAL

A polynomial $p \in \mathbb{R}_{\geq 0}[x_1, \dots, x_n]$ is **log-concave** (at \mathbf{x}) if the Hessian $\nabla^2 \log p(\mathbf{x})$ is negative semi-definite.

$\Rightarrow \nabla^2 p(\mathbf{x})$ has at most one positive eigenvalue.

STRONGLY LOG-CONCAVE POLYNOMIAL

A polynomial $p \in \mathbb{R}_{\geq 0}[x_1, \dots, x_n]$ is **strongly log-concave** if for any index set $I \subseteq [n]$, $\partial_I p$ is log-concave at $\mathbf{1}$.

Originally introduced by [Gurvits \(2009\)](#), equivalent to:

- Completely log-concave ([Anari, Oveis Gharan, and Vintzant, 2018](#));
- Lorentzian polynomials ([Brändén and Huh, 2019+](#)).

STRONGLY LOG-CONCAVE DISTRIBUTIONS

A distribution $\pi : 2^{[n]} \rightarrow \mathbb{R}_{\geq 0}$ is **strongly log-concave** if so is its generating polynomial

$$g_{\pi}(\mathbf{x}) = \sum_{S \subseteq [n]} \pi(S) \prod_{i \in S} x_i.$$

An important example of homogeneous strongly log-concave distributions is the uniform distribution over bases of a matroid (Anari, Oveis Gharan, and Vintzant 2018; Brändén and Huh 2019+).

ALTERNATIVE CHARACTERIZATION FOR SLC

Brändén and Huh (2019+): An r -homogeneous multiaffine polynomial p with non-negative coefficients is **strongly log-concave** if and only if:

- the support of p is the bases of some matroid;
- after taking $r - 2$ partial derivatives, the quadratic is **real stable** or **0**.

Real stable: $p(\mathbf{x}) \neq 0$ if $\Im(x_i) > 0$ for all i .

Real stable polynomials (and strongly Rayleigh distributions) capture only “balanced” matroids, whereas SLC polynomials capture all matroids.

BASES-EXCHANGE WALK

The following Markov chain $P_{\text{BX},\pi}$ converges to a homogeneous SLC π :

- ① **remove** an element uniformly at random from the current basis (call the resulting set S);
- ② **add** $i \notin S$ with probability proportional to $\pi(S \cup \{i\})$.

The implementation of the second step may be non-trivial.

Recall the mixing time defn.

$$t_{\text{mix}}(P, \varepsilon) := \min_t \{t \mid \|P^t(x_0, \cdot) - \pi\|_{\text{TV}} \leq \varepsilon\}.$$

THEOREM (CRYAN-GUO-MOUSA 2019)

For any r -homogeneous strongly log-concave distribution π ,

$$t_{\text{mix}}(P_{\text{BX},\pi}, \varepsilon) \leq r \left(\log \log \frac{1}{\pi_{\min}} + \log \frac{1}{2\varepsilon^2} \right),$$

where $\pi_{\min} = \min_{x \in \Omega} \pi(x)$.

Previously, Anari, Liu, Oveis Gharan, and Vintzant (2019):

$$t_{\text{mix}}(P_{\text{BX},\pi}, \varepsilon) \leq r \left(\log \frac{1}{\pi_{\min}} + \log \frac{1}{\varepsilon} \right)$$

E.g. for the uniform distribution over bases of matroids (with n elements and rank r), the new bound is $O(r(\log r + \log \log n))$, whereas the previous bound is $O(r^2 \log n)$.

LEVELS OF INDEPENDENT SETS

The set of all independent sets of a matroid \mathcal{M} is **downward closed**.

Let $\mathcal{M}(k)$ be the set of independent sets of size k . Thus, $\mathcal{M}(r)$ is the set of all bases.

Let \mathcal{M}_i denote the matroid \mathcal{M} after contracting i , which is another matroid itself.

WEIGHTS FOR INDEPENDENT SETS

Equip \mathcal{M} with the following inductively defined weight function:

$$w(I) := \begin{cases} \pi(I)Z_r & \text{if } |I| = r, \\ \sum_{I' \supset I, |I'|=|I|+1} w(I') & \text{if } |I| < r, \end{cases}$$

for some normalization constant $Z_r > 0$.

For example, we may choose $w(B) = 1$ for all $B \in \mathcal{B}$ and $Z_r = |\mathcal{B}|$, which corresponds to the uniform distribution over \mathcal{B} .

Let π_k be the distribution such that $\pi_k(I) \propto w(I)$, and Z_k be the corresponding normalizing constant.

DIFFERENT VIEWS

- **Polynomial**: $\frac{\partial \mathbf{p}}{\partial x_i}$; setting $x_i = 0$; $(r - k)! \frac{\partial}{\partial x_i} \mathbf{p}(\mathbf{1})$
- **Matroid**: contraction over i ; deletion of i ; $w(I)$
- **Distribution**: conditioning on having i ; conditioning on *not* having i ; proportional to $\pi_k(I)$

RANDOM WALK BETWEEN LEVELS

There are two natural random walks converging to π_k .

The “down-up” random walk P_k^\vee :

- 1. remove an element of $I \in \mathcal{M}(k)$ uniformly at random to get $I' \in \mathcal{M}(k-1)$;
2. move to J such that $J \in \mathcal{M}(k)$, $J \supset I'$ with probability $\frac{w(J)}{w(I')}$.

The bases-exchange walk $P_{\text{BX},\pi} = P_r^\vee$.

The “up-down” walk P_k^\wedge is defined similarly.

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DECOMPOSING THE WALKS

Let A_k be the matrix whose rows are indexed by $\mathcal{M}(k)$ and columns by $\mathcal{M}(k+1)$ such that $A_k(I, J) = 1$ if and only if $I \subset J$.

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Let $\mathbf{w}_k = \{w(I)\}_{I \in \mathcal{M}(k)}$, and

$$P_{k+1}^\downarrow := \frac{1}{k+1} \cdot A_k^\top;$$

$$P_k^\uparrow := \text{diag}(\mathbf{w}_k)^{-1} A_k \text{diag}(\mathbf{w}_{k+1}).$$

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KEY LEMMA

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For any $k \geq 2$ and $f : \mathcal{M}(k) \rightarrow \mathbb{R}_{\geq 0}$,

$$\frac{\text{Ent}_{\pi_k}(f)}{k} \geq \frac{\text{Ent}_{\pi_{k-1}}(P_{k-1}^\uparrow f)}{k-1}.$$

KEY LEMMA

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- Applying P_{k-1}^\uparrow to the left corresponds to the random walk P_k^\downarrow .
- The lemma asserts that P_k^\downarrow contracts the relative entropy by at least $(1 - 1/k)$:

$$\text{Ent}_{\pi_{k-1}}(P_{k-1}^\uparrow f) \leq (1 - 1/k) \text{Ent}_{\pi_k}(f).$$

BASE CASE

For the base case, we want to show that

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Using $a \log \frac{a}{b} \geq a - b$ for $a, b > 0$, we can get

$$\text{Ent}_{\pi_2}(f) - 2\text{Ent}_{\pi_1}(P_1^\uparrow f) \geq 1 - \frac{1}{2Z_2} \cdot h^\top W h,$$

where $W_{ij} = w(\{i, j\})$ and $h = P_1^\uparrow f$.

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where $W_{ij} = w(\{i, j\})$ and $h = P_1^\uparrow f$.

Since $W = (r-2)!Z_r \nabla^2 g_\pi(\mathbf{1})$, it has at most one positive eigenvalue. The quadratic form is maximized at $h = P_1^\uparrow f = \mathbf{1}$, which helps prove the base case...

BOUND THE MIXING TIME DIRECTLY

For a distribution τ on $\mathcal{M}(k)$, the relative entropy $D(\tau \parallel \pi_k) = \text{Ent}_{\pi_k}(D_k^{-1}\tau)$ where $D_k = \text{diag}(\pi_k)$. Moreover, after one step of P_k^\vee , the distribution is $(\tau^T P_k^\vee)^T = (P_k^\vee)^T \tau$. Since P_k^\vee is reversible, $D_k^{-1}(P_k^\vee)^T = P_k^\vee D_k^{-1}$.

$$\begin{aligned}
 D((P_k^\vee)^T \tau \parallel \pi_k) &= \text{Ent}_{\pi_k}(D_k^{-1}(P_k^\vee)^T \tau) \\
 &= \text{Ent}_{\pi_k}(P_k^\vee D_k^{-1} \tau) \\
 &= \text{Ent}_{\pi_k}(P_k^\downarrow P_{k-1}^\uparrow D_k^{-1} \tau) \\
 &\leq \text{Ent}_{\pi_{k-1}}(P_{k-1}^\uparrow D_k^{-1} \tau) && \text{(Jensen's inequality)} \\
 &\leq \left(1 - \frac{1}{k}\right) \text{Ent}_{\pi_k}(D_k^{-1} \tau) && \text{(entropy contraction)} \\
 &= \left(1 - \frac{1}{k}\right) D(\tau \parallel \pi_k).
 \end{aligned}$$

The mixing time bound follows from Pinsker's inequality

$$2 \|\tau - \sigma\|_{\text{TV}}^2 \leq D(\tau \parallel \sigma).$$

MANY OPEN PROBLEMS REMAIN!

- Sampling spanning trees of G with **no broken circuit** - the same as evaluation $T_G(1, 0)$ of the Tutte polynomial, leading coefficient $|a_1[\chi_G(\lambda)]|$ of the chromatic polynomial (up to sign), number of maximum G -parking functions, sampling acyclic orientations *with a unique sink* ...
- Sampling all acyclic orientations of G
- Negative correlation for random forests of G
- (More generally) characterization of matroids with negative correlation for random bases.

FIN

Thank you!