

COMPUTATIONAL REPRESENTATION THEORY – LECTURE III

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Group Theory and Computational Methods
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- 1 Representations of Groups and Algebras
- 2 The MeatAxe

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FG is a finite-dimensional F -algebra.

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From now on: representation = linear representation

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$\mathfrak{x} : G \rightarrow \mathrm{GL}(V)$ and $\mathfrak{y} : G \rightarrow \mathrm{GL}(W)$ are equivalent, if and only if V and W are *isomorphic* as FG -modules.

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Irreducibility and equivalence are defined analogously to the case of group representations.

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- 1 from permutation representations,
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- 4 in various other ways.

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Replacing each $\kappa(g) \in S_\Omega$ by the linear map $\mathfrak{X}(g)$ of V , which permutes its basis like $\kappa(g)$, we obtain an F -representation \mathfrak{X} of degree n of G .

PERMUTATION REPRESENTATIONS: EXAMPLE

Let $G = S_3$, the symmetric group on three letters, acting (from the right) on $\Omega = \{1, 2, 3\}$. Then $G = \langle a, b \rangle$ with $a = (1, 2)$ and $b = (1, 2, 3)$.

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Take $F = \mathbb{Q}$; $V \cong \mathbb{Q}^{1 \times 3}$. Then G acts on V by $e_i \cdot g = e_{ig}$, where e_1, e_2, e_3 is the standard basis of V .

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Then

$$\mathfrak{X}(a) = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad \mathfrak{X}(b) = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

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Note: In order to describe (store) a representation \mathfrak{X} of G , it suffices to give the matrices $\mathfrak{X}(a_i)$ for a generating set $\{a_1, \dots, a_l\}$ of G .

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Here, $\mathfrak{X}(g) \otimes \mathfrak{Y}(g)$ is the **Kronecker product** of the two matrices $\mathfrak{X}(g)$ and $\mathfrak{Y}(g)$, defined as follows:

$$A \otimes B := \begin{bmatrix} a_{11} B & \cdots & a_{1d} B \\ \vdots & \ddots & \vdots \\ a_{d1} B & \cdots & a_{dd} B \end{bmatrix}$$

for $A = [a_{ij}] \in F^{d \times d}$, and $B \in F^{e \times e}$.

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From this: $\chi_{\mathfrak{X} \otimes \mathfrak{Y}} = \chi_{\mathfrak{X}} \cdot \chi_{\mathfrak{Y}}$, i.e. the product of characters is a character.

Let $\mathfrak{X} : G \rightarrow \text{GL}(V)$ be a representation of G on V .

Let W be a G -invariant subspace of V , i.e.:

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We obtain F -representations

$$\mathfrak{X}_W : G \rightarrow \text{GL}(W) \quad \text{and} \quad \mathfrak{X}_{V/W} : G \rightarrow \text{GL}(V/W)$$

in the natural way.

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Choosing the basis $e_1 + e_2 + e_3, e_2, e_3$ of V , and transforming the matrices accordingly, we obtain a representation \mathfrak{X}' with

$$\mathfrak{X}'(a) = \left[\begin{array}{c|cc} 1 & 0 & 0 \\ \hline 1 & -1 & -1 \\ 0 & 0 & 1 \end{array} \right], \quad \mathfrak{X}'(b) = \left[\begin{array}{c|cc} 1 & 0 & 0 \\ \hline 0 & 0 & 1 \\ 1 & -1 & -1 \end{array} \right].$$

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Moreover,

$$\mathfrak{X}'_{V/W}(a) = \begin{bmatrix} -1 & -1 \\ 0 & 1 \end{bmatrix}, \quad \mathfrak{X}'_{V/W}(b) = \begin{bmatrix} 0 & 1 \\ -1 & -1 \end{bmatrix}.$$

THE IRREDUCIBLE CONSTITUENTS

Iterating this process, we obtain a matrix representation \mathfrak{X}^∞ of G , equivalent to \mathfrak{X} , s.t.:

$$\mathfrak{X}^\infty(g) = \begin{bmatrix} \mathfrak{X}_1(g) & 0 & \cdots & 0 \\ * & \mathfrak{X}_2(g) & \cdots & 0 \\ * & * & \ddots & 0 \\ * & * & \cdots & \mathfrak{X}_l(g) \end{bmatrix},$$

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They are unique up to equivalence (Jordan-Hölder theorem).

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one obtains all irreducible F -representations of G .

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Since then, it has been improved and enhanced by many people, including Derek Holt, Gábor Ivanyos, Klaus Lux, Jürgen Müller, Sarah Rees, Michael Ringe, and by Richard Parker himself.

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- 2 How does one prove that \mathfrak{X} is irreducible?

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- 4 \mathfrak{A} acts irreducibly on $F^{1 \times d}$.

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Note: If $G = \langle g_1, \dots, g_l \rangle$, and $\mathfrak{X} : G \rightarrow \mathrm{GL}_d(F)$ is an F -representation of G , take $A_i := \mathfrak{X}(g_i)$, $1 \leq i \leq l$.

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The space $w.\mathfrak{A}$ is computed with a linearized version of an orbit algorithm: $\text{spin}(A_1, \dots, A_l, w)$ returns F -basis of $w.\mathfrak{A}$.

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$\mathcal{B} \leftarrow w$;

for v in \mathcal{B} **do**

for i from 1 to l **do**

if $\{vA_i\} \cup \mathcal{B}$ *linearly independent* **then**

append vA_i to \mathcal{B} ;

end if;

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Holt and Rees use characteristic polynomials of elements of \mathfrak{A} to find suitable B s and also to reduce the number of tests considerably.

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 - if p is good, Return("Irreducible") fi;
- 5 Go back to Step 1.

A huge collection of explicit representations of finite groups is contained in Rob Wilson's *WWW Atlas of Finite Group Representations*: (<http://brauer.maths.qmul.ac.uk/Atlas/>).

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Much of this information is also available through the GAP-package atlasrep (<http://www.math.rwth-aachen.de/~Thomas.Breuer/atlasrep/>).

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Thank you for your attention!